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COMBINATORIAL OPERADS FROM MONOIDS

SAMUELE GIRAUDO

ABSTRACT. We introduce a functorial construction which, from a monoid, produces a set-operad. We obtain new (symmetric or not) operads as suboperads or quotients of the operads obtained from usual monoids such as the additive and multiplicative monoids of integers and cyclic monoids. They involve various familiar combinatorial objects: endofunctions, parking functions, packed words, permutations, planar rooted trees, trees with a fixed arity, Schröder trees, Motzkin words, integer compositions, directed animals, and segmented integer compositions. We also recover some already known (symmetric or not) operads: the magmatic operad, the associative commutative operad, the diassociative operad, and the triassociative operad. We provide presentations by generators and relations of all constructed nonsymmetric operads.

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INTRODUCTION

Operads are algebraic structures introduced in the 1970s by Boardman and Vogt [BV73] and by May [May72] in the context of algebraic topology. Informally, an operad is a structure containing operators with n inputs and 1 output, for all positive integer n . Two operators x and y can be composed at the i th position by grafting the output of y on the i th input of x . The new operator thus obtained is denoted by $x \circ_i y$. In an operad, one can also switch the inputs

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of an operator x by letting a permutation σ act to obtain a new operator denoted by $x \cdot \sigma$. One of the main relishes of operads comes from the fact that they offer a general theory to study in an unifying way different types of algebras, such as associative algebras and Lie algebras.

In recent years, the importance of operads in combinatorics has increased and several new (nonsymmetric) operads were defined on combinatorial objects (see *e.g.*, [Lod01, CL01, Liv06, Cha08]). The structure thereby added on combinatorial families enables us to see these in a new light and offers original ways to solve some combinatorial problems. For example, the dendriform operad [Lod01] is a nonsymmetric operad on binary trees that plays an interesting role in describing the Hopf algebra of Loday-Ronco of binary trees [LR98, HNT05]. Besides, this nonsymmetric operad is a key ingredient for the enumeration of intervals of the Tamari lattice [Cha06, Cha08]. There is also a very rich link connecting combinatorial Hopf algebra theory and operad theory: various constructions produce combinatorial Hopf algebras from operads [vdL04, CL07, Fra08, LV12, ML13].

In this paper, we propose a new generic method to build combinatorial operads. The starting point is to pick a monoid M . We then consider the set of words whose letters are elements of M . The arity of such words are their length, the composition of two words is expressed from the product of M , and permutations act on words by permuting letters. In this way, we associate an operad denoted by $\mathsf{T}M$ with any monoid M . This construction is rich from a combinatorial point of view since it allows us, by considering suboperads and quotients of $\mathsf{T}M$, to get new (symmetric or not) operads on various combinatorial objects. Our construction is related to two previous ones.

The first one is a construction of Méndez and Nava [MN93] emerging from the context of the species theory [Joy81]. Roughly speaking, a species is a combinatorial construction U which takes an underlying set E as input and produces a set $U[E]$ of objects by adding some structure on the elements of E (see [BLL94]). This theory has many links with the theory of operads since an operad is a monoid with respect to the operation of substitution of species. In [MN93], the authors defined the plethystic species, that are species taking as input sets where any element has a colour picked from a fixed monoid M . This monoid has to satisfy some precise conditions (as to be left cancellable and without proper divisor of the unity, and such that any element has finitely many factorizations). It appears that the elements of the so-called uniform plethystic species can be seen as words of colours and hence, as elements of $\mathsf{T}M$. Moreover, the composition of this operad is the one of $\mathsf{T}M$. The main difference between the construction of Méndez and Nava and ours lies in the fact that the construction T can be applied on any monoid.

The second one, introduced by Berger and Moerdijk [BM03], is a construction which allows to obtain, from a commutative bialgebra \mathcal{B} , a cooperad $T\mathcal{B}$. Our construction T and the construction T of these two authors are different but coincide in many cases. For instance, when (M, \bullet) is a monoid such that for any $x \in M$, the set of pairs $(y, z) \in M^2$ satisfying $y \bullet z = x$ is finite, the operad $\mathsf{T}M$ is the dual of the cooperad $T\mathcal{B}$ where \mathcal{B} is the dual bialgebra of $\mathbb{K}[M]$ endowed with the diagonal coproduct (\mathbb{K} is a field). On the other hand, there are operads that we can build by the construction T but not by the construction T , and conversely. For

example, the operad $\mathsf{T}\mathbb{Z}$, where \mathbb{Z} is the additive monoid of integers, cannot be obtained as the dual of a cooperad built by the construction T of Berger and Moerdijk.

Furthermore, our construction is defined in the category of sets and computations are explicit. It is therefore possible given a monoid M , to make experiments on the operad $\mathsf{T}M$, using if necessary a computer. In this paper, we study many applications of the construction T focusing on its combinatorial aspect. More precisely, we define, by starting from very simple monoids like the additive or multiplicative monoids of integers, or cyclic monoids, various nonsymmetric operads involving well-known combinatorial objects.

This paper is organized as follows. In Section 1, we set some notations about syntax trees and rewriting systems on syntax trees. We then briefly recall the basics about operads. We also prove in this section two important lemmas used in the sequel of the paper: the first one deals with the form of elements of nonsymmetric operads generated by a set of generators and the second one is a tool to prove presentations by generators and relations of nonsymmetric operads using rewrite rules on syntax trees. Section 2 defines the construction T , associating an operad with a monoid and establishes its first properties. We show that this construction is a functor from the category of monoids to the category of operads which respects injections and surjections. Finally we apply this construction in Section 3 on various monoids and obtain several new (symmetric or not) operads. We construct in this way some operads on combinatorial objects which were not provided with such a structure: planar rooted trees with a fixed arity, Motzkin words, integer compositions, directed animals, and segmented integer compositions. We also obtain new operads on objects which are already provided with such a structure: endofunctions, parking functions, packed words, permutations, planar rooted trees, and Schröder trees. By using the construction T , we also give an alternative construction for the diassociative operad [Lod01] and for the triassociative operad [LR04].

This paper is an extended version of [Gir12a] and [Gir12b]. It contains all proofs and new results like the presentations by generators and relations of the considered nonsymmetric operads.

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1. SYNTAX TREES AND OPERADS

In this section, we set some notations about syntax trees and operads. All our operads are reduced, *i.e.*, all of its elements have at least one input. In the same way, all our syntax trees are reduced, *i.e.*, all of its internal nodes have at least one child. We also present in this section some notions about rewriting systems on syntax trees that we shall use to prove presentations of operads throughout the rest of the paper.

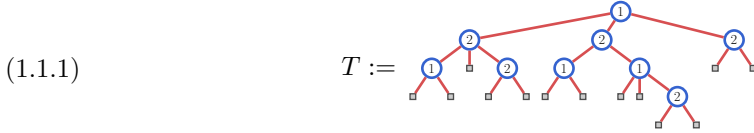
1.1. Rewriting systems on syntax trees. In the same way as the elements of free monoids can be seen as words and also are good objects to study monoids, the elements of free operads can be seen as syntax trees are useful objects to manipulate these algebraic structures.

1.1.1. Syntax trees. We use in the sequel the standard terminology (*i.e.*, *node*, *edge*, *root*, *subtree*, *parent*, *child*, *path*, *etc.*) about trees (see for instance [AU94]). In our graphical representations, internal nodes are represented by circles \circ , leaves by squares \square , and edges by segments ! . Each tree is depicted so that its root is the uppermost node.

Let S be a nonempty set. A *syntax tree on S* , or simply a *syntax tree* if S is fixed, is a planar rooted tree such that internal nodes are labeled on S . We shall denote by \mathcal{T}_n^S the set of syntax trees on S with n leaves and by \mathcal{T}^S the set $\sqcup_{n \geq 1} \mathcal{T}_n^S$.

Let T be a syntax tree. We denote by $n(T)$ (resp. $\ell(T)$) the number of internal nodes (resp. leaves) of T . The *arity* of an internal node is the number of its children. The *depth-first traversal* of T consists in visiting the root of T and then, recursively visiting by a depth-first traversal the subtrees of T , from left to right. The *i th internal node* (resp. *i th leaf*) of T is the i th internal node (resp. i th leaf) of T visited by a depth-first traversal. The *depth* of a node x of T is the length of the unique path connecting x with the root of T . Note that the depth of the root of T is 0. The *weight* $w(T)$ of T is the sum, for all internal nodes x of T , of the number of internal nodes of the rightmost subtree of x .

For example,



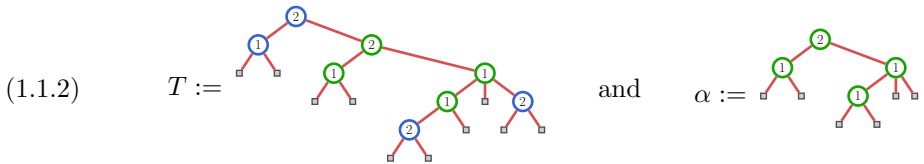
is a syntax tree on the set $\{1, 2\}$. It has 9 internal nodes, 13 leaves, its weight is 5, and the sequence of the labels of its internal nodes visited by the depth-first traversal is 121221122.

1.1.2. Syntax tree patterns. Let T and α be two syntax trees. We say that T *admits an occurrence of α at the root* if one of the following two conditions is satisfied:

- (i) the tree α consists in exactly one leaf and no internal node;
- (ii) the roots of T and α have same arities and same labels, and, by denoting by (T_1, \dots, T_k) and $(\alpha_1, \dots, \alpha_k)$ the sequence of the subtrees of T and α from left to right, T_i admits an occurrence of α_i at the root, for any $i \in [k]$.

We say that T *admits an occurrence of α* if there is a node x of T such that the syntax tree rooted on x admits an occurrence of α at the root.

For example, set



Then T admits one occurrence of α .

1.1.3. Rewrite rules on syntax trees. Let S be a nonempty set and \mathcal{T} be a subset of \mathcal{T}^S . A *rewrite rule* on \mathcal{T} is a binary relation \mapsto on \mathcal{T} such that

$$(1.1.3) \quad \alpha \mapsto \beta \quad \text{implies} \quad \ell(\alpha) = \ell(\beta).$$

A syntax tree T_0 is *rewritable* into a syntax tree T_1 by \mapsto if:

- (i) there are two syntax trees α and β such that $\alpha \mapsto \beta$;
- (ii) there is in T_0 an occurrence of α ;
- (iii) we obtain T_1 by replacing an occurrence of α in T_0 by β .

We denote this property by $T_0 \rightarrow T_1$ and we call the pair (T_0, T_1) a *rewriting*. Moreover, if there is a sequence S_1, \dots, S_k of syntax trees such that

$$(1.1.4) \quad T_0 \rightarrow S_1 \rightarrow \dots \rightarrow S_k \rightarrow T_1,$$

we say that T_0 is *rewritable* into T_1 by \rightarrow and we denote this property by $T_0 \xrightarrow{*} T_1$.

If there is no infinite chain

$$(1.1.5) \quad T_0 \rightarrow T_1 \rightarrow \dots,$$

we say that \mapsto is *terminating*. In this case, a syntax tree T that cannot be rewritten is a *normal form* for \mapsto . When \mapsto is terminating and there are for all $n \geq 1$ finitely many normal forms for \mapsto with n leaves, we denote by $\#(\mapsto_n)$ the number of normal forms for \mapsto with n leaves. In this case, the *generating series* of \mapsto is

$$(1.1.6) \quad F_{\mapsto}(t) := \sum_{n \geq 1} \#(\mapsto_n) t^n.$$

In the sequel, we shall make use of *regular specifications* to describe normal forms of rewrite rules and obtain their generating series. Regular specifications are formal grammars explaining how to build combinatorial objects (see [FS09] for an introduction on regular specifications).

To give an example, set $S := \{1, 2, 3\}$ and consider the rewrite rule \mapsto on \mathcal{T}^S defined by

$$(1.1.7) \quad \begin{array}{c} \text{1} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array} \mapsto \begin{array}{c} \text{3} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array} \quad \text{and} \quad \begin{array}{c} \text{1} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array} \mapsto \begin{array}{c} \text{2} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array}.$$

Here is a sequence of rewritings:

$$(1.1.8) \quad \begin{array}{c} \text{1} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array} \rightarrow \begin{array}{c} \text{3} \\ \diagup \quad \diagdown \\ \text{1} \quad \text{1} \end{array} \rightarrow \begin{array}{c} \text{3} \\ \diagup \quad \diagdown \\ \text{2} \quad \text{1} \end{array} \rightarrow \begin{array}{c} \text{3} \\ \diagup \quad \diagdown \\ \text{2} \quad \text{2} \end{array}.$$

Since for any rewriting $T_0 \rightarrow T_1$, T_1 has less nodes labeled by 1 than T_0 , \mapsto is terminating. The normal forms of \mapsto are the syntax trees on S with no internal node labeled by 1 that has two children.

1.2. Operads. Let us now set, in our context, some definitions and notations about operads. We shall use in the next sections the previous notions about syntax trees to handle elements of free nonsymmetric operads and establish presentations by generators and relations of nonsymmetric operads.

1.2.1. Nonsymmetric operads. A *nonsymmetric operad*, or a *ns operad* for short, is a collection

$$(1.2.1) \quad \mathcal{P} := \bigsqcup_{n \geq 1} \mathcal{P}(n),$$

together with *partial composition maps*

$$(1.2.2) \quad \circ_i : \mathcal{P}(n) \times \mathcal{P}(m) \rightarrow \mathcal{P}(n+m-1), \quad n, m \geq 1, i \in [n],$$

and a distinguished element $\mathbf{1} \in \mathcal{P}(1)$, the *unit* of \mathcal{P} . The above data has to satisfy the following relations:

$$(1.2.3) \quad (x \circ_i y) \circ_{i+j-1} z = x \circ_i (y \circ_j z), \quad x \in \mathcal{P}(n), y \in \mathcal{P}(m), z \in \mathcal{P}(k), i \in [n], j \in [m],$$

$$(1.2.4) \quad (x \circ_i y) \circ_{j+m-1} z = (x \circ_j z) \circ_i y, \quad x \in \mathcal{P}(n), y \in \mathcal{P}(m), z \in \mathcal{P}(k), 1 \leq i < j \leq n,$$

$$(1.2.5) \quad \mathbf{1} \circ_1 x = x = x \circ_i \mathbf{1}, \quad x \in \mathcal{P}(n), i \in [n].$$

One of the simplest ns operads is the *associative commutative operad* \mathbf{Com} . It is defined for all $n \geq 1$ by

$$(1.2.6) \quad \mathbf{Com}(n) := \{\alpha_n\},$$

and the partial composition maps are defined by

$$(1.2.7) \quad \alpha_n \circ_i \alpha_m := \alpha_{n+m-1},$$

for all $n, m \geq 1$ and $i \in [n]$.

Let us now fix some terminology about ns operads and recall basic definitions. The *arity* of an element x of $\mathcal{P}(n)$ is n and is denoted by $|x|$. If \mathcal{Q} is a ns operad, a map $\phi : \mathcal{P} \rightarrow \mathcal{Q}$ is a *ns operad morphism* if it maps elements of arity n of \mathcal{P} to elements of arity n of \mathcal{Q} and commutes with partial composition maps. The ns operad \mathcal{Q} is a *ns suboperad* of \mathcal{P} if $\mathcal{Q} \subseteq \mathcal{P}$, $\mathbf{1} \in \mathcal{Q}$, and \mathcal{Q} is closed for the partial composition maps.

For any set $G \subseteq \mathcal{P}$, the *ns operad generated by G* is the smallest ns suboperad of \mathcal{P} which contains every element of G . When the ns operad generated by G is \mathcal{P} itself and G is minimal with respect to inclusion among the subsets of \mathcal{P} satisfying this property, G is a *generating set* of \mathcal{P} and its elements are *generators* of \mathcal{P} . We say that \mathcal{P} is *finitely generated* if it admits a finite generating set.

The *Hilbert series* of a ns operad \mathcal{P} containing, for all $n \geq 1$, finitely many elements of arity n is the series

$$(1.2.8) \quad F_{\mathcal{P}}(t) := \sum_{n \geq 1} \#\mathcal{P}(n) t^n.$$

A *combinatorial ns operad* is a ns operad which admits an Hilbert series and such that its only element of arity 1 is its unit.

The *composition map* of \mathcal{P} is the mapping

$$(1.2.9) \quad \circ : \mathcal{P}(n) \times \mathcal{P}(m_1) \times \cdots \times \mathcal{P}(m_n) \rightarrow \mathcal{P}(m_1 + \cdots + m_n),$$

defined using partial composition maps \circ_i by

$$(1.2.10) \quad x \circ [y_1, \dots, y_n] := (\dots ((x \circ_n y_n) \circ_{n-1} y_{n-1}) \dots) \circ_1 y_1.$$

The ns operad \mathcal{P} is *basic* if for all $y_1, \dots, y_n \in \mathcal{P}$, the maps

$$(1.2.11) \quad \gamma_{y_1, \dots, y_n} : \mathcal{P}(n) \rightarrow \mathcal{P}(|y_1| + \cdots + |y_n|),$$

defined by

$$(1.2.12) \quad \gamma_{y_1, \dots, y_n}(x) := x \circ [y_1, \dots, y_n]$$

are injective.

1.2.2. Symmetric operads. Let \mathfrak{S}_n be the group of permutations of $[n]$. Any permutation σ is denoted as a word $\sigma_1 \dots \sigma_n$ in such a way that the i th letter σ_i is the image of i . For instance, the word 312 represents the bijection $1 \mapsto 3, 2 \mapsto 1, 3 \mapsto 2$.

To define what is an operad, we need the following definition. Let \mathbf{As} be the ns operad satisfying for all $n \geq 1$,

$$(1.2.13) \quad \mathbf{As}(n) := \mathfrak{S}_n,$$

and, for all $\sigma \in \mathbf{As}(n)$, $\nu \in \mathbf{As}(m)$, and $i \in [n]$,

$$(1.2.14) \quad \sigma \circ_i \nu := \sigma'_1 \dots \sigma'_{i-1} \nu'_1 \dots \nu'_m \sigma'_{i+1} \dots \sigma'_n,$$

where

$$(1.2.15) \quad \nu'_j := \nu_j + \sigma_i - 1, \quad \text{for all } j \in [m],$$

and

$$(1.2.16) \quad \sigma'_j := \begin{cases} \sigma_j & \text{if } \sigma_j < \sigma_i, \\ \sigma_j + m - 1 & \text{otherwise,} \end{cases} \quad \text{for all } j \in [n].$$

This ns operad is known as the *associative noncommutative operad*.

For instance, here are two examples of compositions in \mathbf{As}

$$(1.2.17) \quad \mathbf{123} \circ_2 \mathbf{12} = \mathbf{1234},$$

$$(1.2.18) \quad \mathbf{7415623} \circ_4 \mathbf{231} = \mathbf{941675823}.$$

A *symmetric operad*, or an *operad* for short, is a ns operad \mathcal{P} together with a map

$$(1.2.19) \quad \cdot : \mathcal{P}(n) \times \mathbf{As}(n) \rightarrow \mathcal{P}(n), \quad n \geq 1,$$

which satisfies the following relation:

(1.2.20)

$$(x \cdot \sigma) \circ_i (y \cdot \nu) = (x \circ_{\sigma_i} y) \cdot (\sigma \circ_i \nu), \quad x \in \mathcal{P}(n), y \in \mathcal{P}(m), \sigma \in \text{As}(n), \nu \in \text{As}(m), i \in [n],$$

in such a way that \cdot also is a symmetric group action. Note that any operad \mathcal{P} is also (and thus can be seen as) a ns operad by forgetting its action of As .

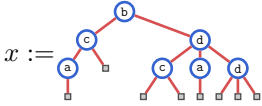
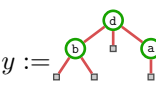
If \mathcal{Q} is an operad, a map $\phi : \mathcal{P} \rightarrow \mathcal{Q}$ is an *operad morphism* if it is a ns operad morphism that commutes with \cdot . The operad \mathcal{Q} is a *suboperad* of \mathcal{P} if \mathcal{Q} is a ns suboperad of \mathcal{P} and \mathcal{Q} is closed for \cdot .

For any set $G \subseteq \mathcal{P}$, the *operad generated by G* is the smallest suboperad of \mathcal{P} which contains every element of G . When the operad generated by G is \mathcal{P} itself and G is minimal with respect to inclusion among the subsets of \mathcal{P} satisfying this property, G is a *generating set* of \mathcal{P} and its elements are *generators* of \mathcal{P} . We say that \mathcal{P} is *finitely generated* if it admits a finite generating set.

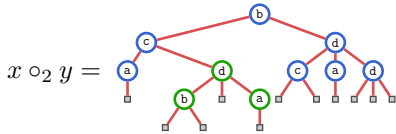
1.3. Presentation of nonsymmetric operads. We now focus on ns operads and present the tools we will need to establish presentations by generators and relations of ns operads.

1.3.1. Free ns operads. Let $S := \sqcup_{n \geq 1} S(n)$ be a set. The *free ns operad* over S is the ns operad $\mathcal{F}(S)$ defined as follows. For any $n \geq 1$, the set $\mathcal{F}(S)(n)$ is the set of syntax trees on S with n leaves and where internal nodes which have k children are labeled on $S(k)$. The composition $x \circ_i y$ of two elements of $\mathcal{F}(S)$ consists in grafting the root of y on the i th leaf of x . The unit $\mathbf{1}$ of $\mathcal{F}(S)$ is the tree with no internal node and hence exactly one leaf. The *degree* $d(x)$ of an element x of $\mathcal{F}(S)$ is its number $n(x)$ of internal nodes and its *arity* $|x|$ is its number $\ell(x)$ of leaves. If $d(x) = 1$, the *label* $\text{lbl}(x)$ of x is the element of S which labels the only internal node of x .

Let $S := S(1) \sqcup S(2) \sqcup S(3)$, where $S(1) := \{\mathbf{a}\}$, $S(2) := \{\mathbf{b}, \mathbf{c}\}$, and $S(3) := \{\mathbf{d}\}$. Then,

(1.3.1)  and 

are two elements of $\mathcal{F}(S)$. The arity of x is 8 and its degree is 7. The arity of y is 4 and its degree is 3. Moreover, one has in $\mathcal{F}(S)$ the following composition

(1.3.2) 

1.3.2. *Evaluations.* Let $(\mathcal{P}, \circ_i^{\mathcal{P}}, \mathbf{1}_{\mathcal{P}})$ be a ns operad and $X := \sqcup_{n \geq 1} X(n)$ be a set of elements of \mathcal{P} , where $X(n) \subseteq \mathcal{P}(n)$ for all $n \geq 1$. The *evaluation* ev is the mapping

$$(1.3.3) \quad \text{ev} : \mathcal{F}(X) \rightarrow \mathcal{P},$$

recursively defined by $\text{ev}(\mathbf{1}) := \mathbf{1}_{\mathcal{P}}$, $\text{ev}(x) := \text{lbl}(x)$ if $d(x) = 1$, and

$$(1.3.4) \quad \text{ev}(x) := \text{ev}(y) \circ_i^{\mathcal{P}} \text{ev}(z)$$

if $d(x) \geq 2$, where $y, z \in \mathcal{F}(X) \setminus \{\mathbf{1}\}$, $i \in [|y|]$, and $x = y \circ_i z$.

In other words, we can see an element x of $\mathcal{F}(X)$ as a tree-like expression for an element $\text{ev}(x)$ of \mathcal{P} . Moreover, it is easy, by induction on the degree, to prove that ev is a well-defined surjective mapping and hence, is a ns operad morphism.

1.3.3. *Ns operadic congruences.* An *ns operadic congruence* over a ns operad \mathcal{P} is an equivalence relation \equiv on its elements such that $x \equiv x'$ implies $|x| = |x'|$, and for all $x, x', y, y' \in \mathcal{P}$ and $i \in [|x|]$,

$$(1.3.5) \quad x \equiv x' \text{ and } y \equiv y' \quad \text{implies} \quad x \circ_i y \equiv x' \circ_i y'.$$

Given a set $S := \sqcup_{n \geq 1} S(n)$ and an operadic congruence \equiv over the free ns operad $\mathcal{F}(S)$, one can construct a *ns quotient operad* $\mathcal{F}(S)/_{\equiv}$ of $\mathcal{F}(S)$ defined as follows. We set

$$(1.3.6) \quad \mathcal{F}(S)/_{\equiv}(n) := \{[x]_{\equiv} : x \in \mathcal{F}(S)(n)\}, \quad n \geq 1,$$

where $[x]_{\equiv}$ is the \equiv -equivalence class of x , and

$$(1.3.7) \quad [x]_{\equiv} \circ_i [y]_{\equiv} := [x \circ_i y]_{\equiv},$$

where x and y are any elements of $\mathcal{F}(S)$ such that $x \in [x]_{\equiv}$ and $y \in [y]_{\equiv}$.

1.3.4. *Presentation by generators and relations.* In the sequel, we shall define ns operadic congruences \equiv over free ns operads $\mathcal{F}(S)$ through equivalence relations \leftrightarrow on the set $\mathcal{F}(S)$. The *congruence generated by* \leftrightarrow is the most refined ns operadic congruence \equiv containing \leftrightarrow .

Besides, we say that a relation \mapsto on $\mathcal{F}(S)$ is an *orientation* of \leftrightarrow if \mapsto is the finest relation such that its reflexive, symmetric, and transitive closure is \leftrightarrow . The link between rewrite rules on syntax trees and ns operads relies on the fact that orientations can be regarded as rewrite rules.

A *presentation* of a ns operad \mathcal{P} consists in a set $S := \sqcup_{n \geq 1} S(n)$ and a ns operadic congruence \equiv over $\mathcal{F}(S)$ such that $\mathcal{P} = \mathcal{F}(S)/_{\equiv}$. When $S(2) \neq \emptyset$ and $S(n) = \emptyset$ for all $n \neq 2$, \mathcal{P} is called *binary*. When \equiv can be generated as a ns operadic congruence by an equivalence relation \leftrightarrow on $\mathcal{F}(S)$ only involving elements of degree 2, \mathcal{P} is called *quadratic*.

The following lemma presents a description of the elements of ns operads generated by a set of generators.

Lemma 1.1. *Let \mathcal{P} be a ns operad generated by a set G of generators. Then any element x of \mathcal{P} different from the unit of \mathcal{P} can be written as*

$$(1.3.8) \quad x = y \circ_i g,$$

where $y \in \mathcal{P}(n)$, $n \geq 1$, $g \in G$, and $i \in [n]$.

Proof. Since the map $\text{ev} : \mathcal{F}(G) \rightarrow \mathcal{P}$ is surjective, x admits a tree-like expression $x' \in \mathcal{F}(G)$ satisfying $\text{ev}(x') = x$. Since x is different from the unit of \mathcal{P} , one has $x' = y' \circ_i g'$ for two syntax trees y' and g' of $\mathcal{F}(G)$ such that g' has exactly one internal node. Then, by setting $y := \text{ev}(y')$ and $g := \text{ev}(g')$, we have, since ev is a ns operad morphism, $x = y \circ_i g$. \square

We shall use in the sequel Lemma 1.1 to study ns operads \mathcal{Q} generated by a subset of elements of a bigger ns operad \mathcal{P} . It allows us to describe \mathcal{Q} arity by arity because any element of \mathcal{Q} can be obtained by composing an element of a smaller arity with a generator.

The following lemma presents a tool for showing that a given combinatorial ns operad admits a specified presentation.

Lemma 1.2. *Let \mathcal{P} be a combinatorial ns operad generated by a set G of generators and \equiv be a ns operadic congruence over $\mathcal{F}(G)$ generated by an equivalence relation \leftrightarrow on $\mathcal{F}(G)$. If the following two conditions are satisfied together:*

- (i) *for all $x, x' \in \mathcal{F}(G)$, $x \leftrightarrow x'$ implies $\text{ev}(x) = \text{ev}(x')$;*
- (ii) *there exists an orientation \mapsto of \leftrightarrow such that \mapsto is terminating and has as many normal forms of arity n as elements of \mathcal{P} of arity n ;*

then \mathcal{P} admits the presentation $\mathcal{P} = \mathcal{F}(G)/\equiv$.

Proof. The definition of the evaluation map ev and (i) imply that the map

$$(1.3.9) \quad \phi : \mathcal{F}(G)/\equiv \rightarrow \mathcal{P}$$

defined for any $x \in \mathcal{F}(G)$ by $\phi([x]_{\equiv}) := \text{ev}(x)$ is a surjective ns operad morphism.

Since \mapsto is an orientation of \leftrightarrow , for any $x \in \mathcal{F}(G)$, there is at least one normal form for \mapsto in $[x]_{\equiv}$ and by (ii), for all $n \geq 1$, we have

$$(1.3.10) \quad \#\mathcal{P}(n) = \#(\mapsto_n) \geq \#\mathcal{F}(G)/\equiv(n).$$

This, together with the fact that ϕ is surjective, implies that ϕ also is an isomorphism. Hence, \mathcal{P} admits the claimed presentation. \square

2. A COMBINATORIAL FUNCTOR FROM MONOIDS TO OPERADS

We describe in this section the main ingredient of this paper, namely the *construction* T . This functorial construction associates an operad $\mathsf{T}M$ with any monoid M and an operad morphism $\mathsf{T}\theta : \mathsf{T}M \rightarrow \mathsf{T}N$ with any monoid morphism $\theta : M \rightarrow N$.

2.1. The construction.

2.1.1. *From monoids to operads.* Let $(M, \bullet, 1)$ be a monoid. Let us denote by TM the collection

$$(2.1.1) \quad \mathsf{TM} := \bigsqcup_{n \geq 1} \mathsf{TM}(n),$$

where for all $n \geq 1$,

$$(2.1.2) \quad \mathsf{TM}(n) := \{(x_1, \dots, x_n) : x_i \in M \text{ for all } i \in [n]\}.$$

We endow the set TM with maps

$$(2.1.3) \quad \circ_i : \mathsf{TM}(n) \times \mathsf{TM}(m) \rightarrow \mathsf{TM}(n + m - 1), \quad n, m \geq 1, i \in [n],$$

defined as follows: for all $x \in \mathsf{TM}(n)$, $y \in \mathsf{TM}(m)$, and $i \in [n]$,

$$(2.1.4) \quad x \circ_i y := (x_1, \dots, x_{i-1}, x_i \bullet y_1, \dots, x_i \bullet y_m, x_{i+1}, \dots, x_n).$$

Let us also set $\mathbf{1} := (1)$ as a distinguished element of $\mathsf{TM}(1)$. We endow finally each set $\mathsf{TM}(n)$ with a map

$$(2.1.5) \quad \cdot : \mathsf{TM}(n) \times \mathsf{As}(n) \rightarrow \mathsf{TM}(n), \quad n \geq 1,$$

defined as follows: for all $x \in \mathsf{TM}(n)$ and $\sigma \in \mathsf{As}(n)$,

$$(2.1.6) \quad x \cdot \sigma := (x_{\sigma_1}, \dots, x_{\sigma_n}).$$

The elements of TM are words over M regarded as an alphabet. The arity $|x|$ of an element x of $\mathsf{TM}(n)$ is n . For the sake of readability, we shall denote in some cases an element (x_1, \dots, x_n) of $\mathsf{TM}(n)$ by its *word notation* $x_1 \dots x_n$.

Proposition 2.1. *If M is a monoid, then TM is an operad.*

Proof. This is a straightforward checking of the relations of operads: (1.2.3) comes from the fact that the product of M is associative, (1.2.4) comes from the fact that the elements of TM are words, (1.2.5) comes from the fact that M has a unit, and (1.2.20) comes from the fact that \cdot acts by permuting the letters of the words. \square

2.1.2. *From monoids morphisms to operads morphisms.* Let M and N be two monoids and $\theta : M \rightarrow N$ be a monoid morphism. Let us denote by $\mathsf{T}\theta$ the map

$$(2.1.7) \quad \mathsf{T}\theta : \mathsf{TM} \rightarrow \mathsf{TN},$$

defined for all $(x_1, \dots, x_n) \in \mathsf{TM}(n)$ by

$$(2.1.8) \quad \mathsf{T}\theta(x_1, \dots, x_n) := (\theta(x_1), \dots, \theta(x_n)).$$

Proposition 2.2. *If M and N are two monoids and $\theta : M \rightarrow N$ is a monoid morphism, then the map $\mathsf{T}\theta : \mathsf{TM} \rightarrow \mathsf{TN}$ is an operad morphism.*

Proof. This is a straightforward checking: the fact that θ is a monoid morphism implies the statement of the proposition. \square

2.2. Main properties of the construction.

2.2.1. Functoriality of T .

Proposition 2.3. *Let M and N be two monoids and $\theta : M \rightarrow N$ be a monoid morphism. If θ is injective (resp. surjective), then $\mathsf{T}\theta$ is injective (resp. surjective).*

Proof. This is a straightforward checking: the fact that $\mathsf{T}\theta$ acts letter by letter implies the statement of the proposition. \square

Theorem 2.4. *The construction T is a functor from the category of monoids with monoid morphisms to the category of operads with operad morphisms. Moreover, T respects injections and surjections.*

Proof. By Proposition 2.1, T constructs an operad from a monoid, and by Proposition 2.2, an operad morphism from a monoid morphism. Now, since T sends identity monoid morphisms to identity operad morphisms and T commutes with map composition, T is a functor. Finally, by Proposition 2.3, T also respects injections and surjections, whence the statement of the theorem. \square

2.2.2. *Miscellaneous properties.* Recall that a monoid (M, \bullet) is *right cancellable* if for any $x, y, z \in M$, $y \bullet x = z \bullet x$ implies $y = z$.

Proposition 2.5. *Let M be a monoid. The operad $\mathsf{T}M$ is basic if and only if M is a right cancellable monoid.*

Proof. Let us denote by \bullet the product of M .

Assume first that M is a right cancellable monoid. Let $y^{(1)}, \dots, y^{(n)} \in \mathsf{T}M$, $x, x' \in \mathsf{T}M(n)$, and assume that

$$(2.2.1) \quad \gamma_{y^{(1)}, \dots, y^{(n)}}(x) = \gamma_{y^{(1)}, \dots, y^{(n)}}(x').$$

Then, for any $i \in [n]$ and $j \in [|y^{(i)}|]$,

$$(2.2.2) \quad x_i \bullet y_j^{(i)} = x'_i \bullet y_j^{(i)}.$$

Since M is right cancellable, $x_i = x'_i$ and then, $x = x'$. This implies that $\gamma_{y^{(1)}, \dots, y^{(n)}}$ is injective and that $\mathsf{T}M$ is basic.

Conversely, assume now that $\mathsf{T}M$ is basic. In particular, for any $y \in \mathsf{T}M(1)$, the map γ_y is injective. Hence, for any $x, x' \in \mathsf{T}M(1)$, the equality $\gamma_y(x) = \gamma_y(x')$ implies $x = x'$. This is equivalent to say that $x \bullet y = x' \bullet y$ implies $x = x'$. This amounts exactly to say that M is a right cancellable monoid. \square

Proposition 2.6. *Let M be a monoid generated by a set G . The ns operad $\mathsf{T}M$ is generated by the set*

$$(2.2.3) \quad \{(g) : g \in G\} \cup \{(1, 1)\},$$

where $(1, 1) \in \mathsf{T}M(2)$ and 1 is the unit of M .

Proof. Any element $x := (x_1, \dots, x_n)$ of TM can be generated by the elements of (2.2.3) in the following way. First, generate the element $y := (1, \dots, 1)$ of arity n by composing $(1, 1)$ with itself $n - 1$ times. Next, change each letter y_i of y by composing y with a sequence of generators of G to reach x_i . This is possible since M is generated by G . \square

Theorem 2.7. *Let (M, \bullet) be a monoid generated by a set $G := \{g_1, g_2, \dots\}$ of generators satisfying a set R of nontrivial relations. Then, any algebra S over the ns operad TM is a set equipped with maps*

$$(2.2.4) \quad \star : S \times S \rightarrow S$$

and

$$(2.2.5) \quad \uparrow_g : S \rightarrow S, \quad g \in G$$

satisfying for all $a, b, c \in S$, $g \in G$, and all relations $g_{i_1} \bullet \dots \bullet g_{i_n} = g_{j_1} \bullet \dots \bullet g_{j_m}$ of R , the equalities

$$(2.2.6) \quad (a \star b) \star c = a \star (b \star c),$$

$$(2.2.7) \quad (a \star b) \uparrow_g = a \uparrow_g \star b \uparrow_g,$$

$$(2.2.8) \quad a \uparrow_{g_{i_1}} \dots \uparrow_{g_{i_n}} = a \uparrow_{g_{j_1}} \dots \uparrow_{g_{j_m}}.$$

Proof. Proving the statement of the theorem is equivalent to prove that the ns operad TM admits the presentation by generators and relations obtained by traducing (2.2.4), (2.2.5), (2.2.6), (2.2.7), and (2.2.8) in operadic terms. Thereby, this ns operad \mathcal{P} is the quotient of the free operad generated by a binary generator \star and unary generators \uparrow_g , $g \in G$, submitted to the relations

$$(2.2.9) \quad \star \circ_1 \star = \star \circ_2 \star,$$

$$(2.2.10) \quad \uparrow_g \circ_1 \star = \star \circ [\uparrow_g, \uparrow_g], \quad g \in G,$$

$$(2.2.11) \quad \uparrow_{g_{i_1}} \circ_1 \dots \circ_1 \uparrow_{g_{i_n}} = \uparrow_{g_{j_1}} \circ_1 \dots \circ_1 \uparrow_{g_{j_m}}$$

for all relations $g_{i_1} \bullet \dots \bullet g_{i_n} = g_{j_1} \bullet \dots \bullet g_{j_m}$ of R .

Let $\phi : \mathcal{P} \rightarrow \mathsf{TM}$ be ns operad morphism defined by $\phi(\star) := (1, 1)$ and $\phi(\uparrow_g) := (g)$ for any $g \in G$, where 1 denotes the unit of M . This morphism is well-defined since the elements $(1, 1)$ and (g) of TM satisfy the above relations by replacing \star by $(1, 1)$ and \uparrow_g by (g) . Proposition 2.6 implies that ϕ is surjective since, as a ns operad, TM is generated by $(1, 1)$ and (g) , $g \in G$.

Now, since the equivalence classes of \mathcal{P} are clearly in bijection with the elements of TM , this shows that ϕ is an isomorphism. \square

Monoid	Ns operad	Generators	First dimensions	Combinatorial objects
\mathbb{N}	End	—	1, 4, 27, 256, 3125	Endofunctions
	PF	—	1, 3, 16, 125, 1296	Parking functions
	PW	—	1, 3, 13, 75, 541	Packed words
	Per	—	1, 2, 6, 24, 120	Permutations
	PRT	01	1, 1, 2, 5, 14, 42	Planar rooted trees
	FCat ^(k)	00, 01, ..., 0k	Fuß-Catalan numbers	k -leafy trees
	Schr	00, 01, 10	1, 3, 11, 45, 197	Schröder trees
	Motz	00, 010	1, 1, 2, 4, 9, 21, 51	Motzkin words
\mathbb{N}_2	Comp	00, 01	1, 2, 4, 8, 16, 32	Int. compo.
\mathbb{N}_3	DA	00, 01	1, 2, 5, 13, 35, 96	Directed animals
	SComp	00, 01, 02	1, 3, 27, 81, 243	Seg. int. compo.

TABLE 1. Ground monoids, generators, first dimensions, and combinatorial objects involved in the ns suboperads and quotients of \mathbb{TN} .

3.1.1. *Operads on endofunctions, parking functions, packed words, and permutations.* Recall that an *endofunction* of size n is a word x of length n on the alphabet $\{1, \dots, n\}$. A *parking function* of size n is an endofunction x of size n such that the nondecreasing rearrangement y of x satisfies $y_i \leq i$ for all $i \in [n]$. A *packed word* of size n is an endofunction x of size n such that for any letter $x_i \geq 2$ of x , there is in x a letter $x_j = x_i - 1$.

Note that neither the set of endofunctions nor the set of parking functions, packed words, and permutations are suboperads of \mathbb{TN} . Indeed, one has the following counterexample:

$$(3.1.1) \quad \mathbf{12} \circ_2 \mathbf{12} = \mathbf{134},$$

and, even if 12 is a permutation, 134 is not an endofunction.

Therefore, let us call a word x a *twisted* endofunction (resp. parking function, packed word, permutation) if the word $(x_1 + 1, x_2 + 1, \dots, x_n + 1)$ is an endofunction (resp. parking function, packed word, permutation). For example, the word 2300 is a twisted endofunction since 3411 is an endofunction. Let us denote by End (resp. PF, PW, Per) the set of twisted endofunctions (resp. parking functions, packed words, permutations). Under this reformulation, one has the following result:

Proposition 3.1. *The sets End, PF, and PW form suboperads of \mathbb{TN} .*

Proof. First, by definition of the partial composition map of \mathbb{TN} , the set of twisted endofunctions forms a suboperad of \mathbb{TN} .

Let x and y be two twisted parking functions (resp. packed words) and $i \in [|x|]$. Since x and y have by definition at least one occurrence of 0, we have in \mathbb{TN} ,

$$(3.1.2) \quad \text{Alph}(x \circ_i y) = \text{Alph}(x) \cup \{x_i + a : a \in \text{Alph}(y)\},$$

where $\text{Alph}(u)$ is the set $\{u_j : j \in [|u|]\}$. This, in addition to the fact that any permutation of a twisted parking function (resp. packed word) is still a twisted parking functions (resp. packed word), shows that the partial composition maps of TN and the map \cdot are still well-defined in PF (resp. PW). \square

For example, we have in End the following composition

$$(3.1.3) \quad 2123 \circ_2 30313 = 24142423,$$

and the following application of the map \cdot

$$(3.1.4) \quad 11210 \cdot 23514 = 12011.$$

Note that End is not a finitely generated operad. Indeed, the twisted endofunctions x of size n satisfying $x_i := n - 1$ for all $i \in [n]$ cannot be obtained by compositions involving elements of End of arity smaller than n . Similarly, PF is not a finitely generated operad since the twisted parking functions x of size n satisfying $x_i := 0$ for all $i \in [n - 1]$ and $x_n := n - 1$ cannot be obtained by compositions involving elements of PF of arity smaller than n .

However, the operad PW is a finitely generated operad:

Proposition 3.2. *The operad PW is the suboperad of TN generated by the elements 00 and 01.*

Proof. Let \mathcal{P} be the suboperad of TN generated by the elements 00 and 01, and let us show that $\mathcal{P} = \text{PW}$.

First, by Proposition 3.1, since 00 and 01 are twisted packed words, the elements of \mathcal{P} also are twisted packed words.

Now let x be a nondecreasing twisted packed word and let us show by induction on the size of x that $x \in \mathcal{P}$. If $|x| = 1$, since x is a twisted packed word, one has $x = 0$ and since 0 is the unit of TN , $x \in \mathcal{P}$. Otherwise, let y be the prefix of size $n - 1$ of x . Since x is a nondecreasing word, there are two possibilities to express the last letter x_n of x from the letter x_{n-1} . If $x_n = x_{n-1}$, we have $x = y \circ_{n-1} 00$, and if $x_n = x_{n-1} + 1$, we have $x = y \circ_{n-1} 01$. Hence, since by induction hypothesis \mathcal{P} contains y , \mathcal{P} also contains x . Finally, since any twisted packed word z can be obtained from a nondecreasing packed word x by permuting its letters, we have $z = x \cdot \sigma$ for a certain permutation σ of $\text{As}(n)$, and hence, $\mathcal{P} = \text{PW}$. \square

Let \mathbb{K} be a field and let us from now consider that PW is an operad in the category of \mathbb{K} -vector spaces, *i.e.*, PW is the free \mathbb{K} -vector space over the set of twisted packed words with partial composition maps and the map \cdot extended by linearity. For more details on operads in the category of vector spaces, we redirect the reader to [LV12].

Let I be the free \mathbb{K} -vector space over the set of twisted packed words having multiple occurrences of a same letter.

Proposition 3.3. *The vector space I is an operadic ideal of PW . Moreover, the operadic quotient PW/I is the free vector space over the set of twisted permutations Per and, for all*

twisted permutations x and y , the partial composition map in Per is expressed as

$$(3.1.5) \quad x \circ_i y = \begin{cases} x \circ_i y & \text{if } x_i = |x|, \\ 0_{\mathbb{K}} & \text{otherwise,} \end{cases}$$

where $0_{\mathbb{K}}$ is the null vector of Per and the partial composition map \circ_i in the right member of (3.1.5) is the partial composition map of PW .

Proof. Let x be a twisted packed word and y be a twisted packed word having multiple occurrences of a same letter. Since x and y have at least one occurrence of 0, any composition involving x and y also has multiple occurrences of a same letter. Moreover, for any permutation σ of As of size $|y|$, $y \cdot \sigma$ also has multiple occurrences of a same letter. Hence, I is an operadic ideal of PW and one can consider the operadic quotient PW/I .

Since twisted packed words with no multiple occurrence of a same letter are twisted permutations, PW/I can be identified with the \mathbb{K} -vector space over the set of twisted permutations Per and (3.1.5) follows from the fact that the composition $x \circ_i y$ of two twisted permutations x and y is still a twisted permutation if and only if x_i is the greatest letter of x . \square

Here are two examples of compositions in Per

$$(3.1.6) \quad \mathbf{20431} \circ_1 \mathbf{102} = 0_{\mathbb{K}},$$

$$(3.1.7) \quad \mathbf{20431} \circ_3 \mathbf{102} = \mathbf{2054631}.$$

3.1.2. *A ns operad on planar rooted trees.* Let PRT be the ns suboperad of TN generated by 01. The following table shows the first elements of PRT .

Arity	Elements of PRT
1	0
2	01
3	011, 012
4	0111, 0112, 0121, 0122, 0123
5	01111, 01112, 01121, 01122, 01123, 01211, 01212, 01221, 01222, 01223, 01231, 01232, 01233, 01234

One has the following characterization of the elements of PRT :

Proposition 3.4. *The elements of PRT are exactly the words x on the alphabet \mathbb{N} satisfying $x_1 = 0$ and $1 \leq x_{i+1} \leq x_i + 1$ for all $i \in [|x| - 1]$.*

Proof. Let us first show by induction on the length of the words that any word x of PRT satisfies the statement. This is true when $|x| = 1$. When $|x| \geq 2$, by Lemma 1.1, there is an element y of PRT of length $n := |x| - 1$ and an integer $i \in [n]$ such that $x = y \circ_i 01$. We have

$$(3.1.8) \quad x = (y_1, \dots, y_{i-1}, y_i, y_i + 1, y_{i+1}, \dots, y_n).$$

Since $x_{i+1} = x_i + 1$ and since, by induction hypothesis, y satisfies the statement, x also satisfies it.

Let us now show by induction on the length of the words that PRT contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, since $x_1 = 0$ and $x_2 = 1$, x has a factor $x_i x_{i+1}$ where i is the greatest integer such that $x_{i+1} = x_i + 1$. Now, by setting

$$(3.1.9) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n),$$

we have $x = y \circ_i 01$, and, since i is maximal, if $i + 2 \leq n$ we have $x_{i+2} \leq x_{i+1}$. This implies that y satisfies the statement. By induction hypothesis, PRT contains y and, since $x = y \circ_i 01$, PRT also contains x . \square

Recall that there are $\frac{1}{n} \binom{2n-2}{n-1}$ planar rooted trees with n nodes. There is a bijection ϕ_{PRT} between the words of PRT of arity n and planar rooted trees with n nodes.

To compute $\phi_{\text{PRT}}(x)$ where x is an element of PRT , iteratively insert the letters of x from left to right according to the following procedure. If $|x| = 1$, then $x = 0$ and $\phi_{\text{PRT}}(0)$ is the only planar rooted tree with one node. Otherwise, the insertion of a letter $a \geq 1$ into a planar rooted tree T consists in grafting in T a new node as the rightmost child of the last node of depth $a - 1$ for the depth-first traversal of T .

The inverse bijection is computed as follows. Given a planar rooted tree T of size n , one computes an element of PRT of arity n by labelling each node of T by its depth and then, by reading its labels following a depth-first traversal of T .

Since the elements of PRT satisfy Proposition 3.4, ϕ_{PRT} is well-defined. Hence, we can regard the elements of arity n of PRT as planar rooted trees with n nodes. Figure 2 shows an example of this bijection.

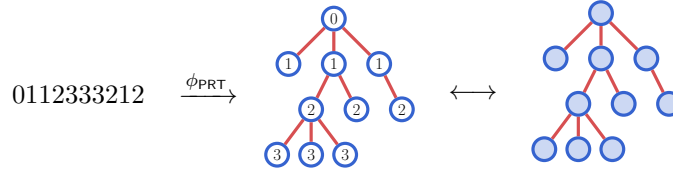


FIGURE 2. Interpretation of an element of the ns operad PRT in terms of planar rooted trees via the bijection ϕ_{PRT} . The nodes of the planar rooted tree in the middle are labeled by their depth.

The bijection ϕ_{PRT} between elements of PRT and planar rooted trees offers an alternative way to compute the composition of elements of PRT :

Proposition 3.5. *Let S and T be two planar rooted trees and s be the i th node for the depth-first traversal of S . The composition $S \circ_i T$ in PRT amounts to replace s by the root of T and graft the children of s as rightmost sons of the root of T .*

Proof. Let $x \in \text{PRT}(n)$ and $y \in \text{PRT}(m)$ such that $S := \phi_{\text{PRT}}(x)$ and $T := \phi_{\text{PRT}}(y)$. Let $U := \phi_{\text{PRT}}(x \circ_i y)$. By definition of ϕ_{PRT} and the partial composition maps of PRT , U is obtained by inserting the prefix of length $i - 1$ of x , then the letters of y incremented by x_i , and finally, the suffix of length $n - i$ of x . Since by Proposition 3.4, y starts by 0, the nodes created by inserting the letters of y incremented by x_i are descendants of the node created by inserting $x_i = y_1$. Moreover, the nodes corresponding to the letters of the suffix of length $n - i$ of x have same parents as they have in T . This implies the statement. \square

Figure 3 shows an example of composition in PRT .

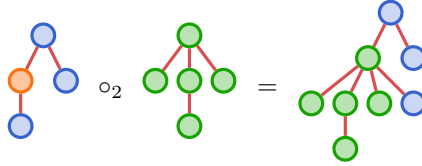


FIGURE 3. Interpretation of the partial composition map of the ns operad PRT in terms of planar rooted trees.

Proposition 3.6. *The ns operad PRT is isomorphic to the free ns operad generated by one element of arity 2.*

Proof. By the characterization of its elements given by Proposition 3.4 and the bijection ϕ_{PRT} , there are as many elements in PRT of arity n than elements of arity n of the free ns operad generated by one element of arity 2. These two ns operads are hence isomorphic. \square

Proposition 3.6 also says that PRT is isomorphic to the magmatic operad and hence, that PRT is a realization of the magmatic operad. This result is already known since in [MY91], Méndez and Yang point out that the species of parenthesizations (binary trees) and the species of planar rooted trees are isomorphic. This isomorphism implies that these species are also isomorphic as ns operads. Moreover, PRT can be seen as a planar version of the *non-associative permutative operad* NAP [MY91] (see also [Liv06]) seen as a ns operad, which is an operad involving labeled non-planar rooted trees.

3.1.3. A ns operad on leafy trees with a fixed arity. Let $k \geq 0$ be an integer and $\text{FCat}^{(k)}$ be the ns suboperad of TN generated by $00, 01, \dots, 0k$. The following tables, respectively, show the first elements of $\text{FCat}^{(1)}$ and $\text{FCat}^{(2)}$.

Arity	Elements of $\text{FCat}^{(1)}$
1	0
2	00, 01
3	000, 001, 010, 011, 012
4	0000, 0001, 0010, 0011, 0012, 0100, 0101, 0110, 0111, 0112, 0120, 0121, 0122, 0123

Arity	Elements of $\mathbf{FCat}^{(2)}$
1	0
2	00, 01, 02
3	000, 001, 002, 010, 011, 012, 013, 020, 021, 022, 023, 024
4	0000, 0001, 0002, 0010, 0011, 0012, 0013, 0020, 0021, 0022, 0023, 0024, 0100, 0101, 0102, 0110, 0111, 0112, 0113, 0120, 0121, 0122, 0123, 0124, 0130, 0131, 0132, 0133, 0134, 0135, 0200, 0201, 0202, 0210, 0211, 0212, 0213, 0220, 0221, 0222, 0223, 0224, 0230, 0231, 0232, 0233, 0234, 0235, 0240, 0241, 0242, 0243, 0244, 0245, 0246

It is immediate from the definition of $\mathbf{FCat}^{(k)}$ that for any $k \geq 0$, $\mathbf{FCat}^{(k)}$ is a ns suboperad of $\mathbf{FCat}^{(k+1)}$. Hence, the ns operads $\mathbf{FCat}^{(k)}$ form an increasing sequence (for inclusion) of ns operads. Note that $\mathbf{FCat}^{(0)}$ is isomorphic to the associative commutative operad \mathbf{Com} . Note also that since $\mathbf{FCat}^{(1)}$ is generated by 00 and 01 and since \mathbf{PRT} is generated by 01, \mathbf{PRT} is a ns suboperad of $\mathbf{FCat}^{(1)}$. Moreover, $\mathbf{FCat}^{(0)}$ is a quotient of \mathbf{PRT} by the ns operadic congruence \equiv defined for all $x, y \in \mathbf{PRT}(n)$ by $x \equiv y$.

One has the following characterization of the elements of $\mathbf{FCat}^{(k)}$:

Proposition 3.7. *The elements of $\mathbf{FCat}^{(k)}$ are exactly the words x on the alphabet \mathbb{N} satisfying $x_1 = 0$ and $0 \leq x_{i+1} \leq x_i + k$ for all $i \in [|x| - 1]$.*

Proof. Let us first show by induction on the length of the words that any word x of $\mathbf{FCat}^{(k)}$ satisfies the statement. This is true when $|x| = 1$. When $|x| \geq 2$, by Lemma 1.1, there is an element y of $\mathbf{FCat}^{(k)}$ of length $n := |x| - 1$, an integer $i \in [n]$, and $0 \leq h \leq k$ such that $x = y \circ_i 0h$. We have

$$(3.1.10) \quad x = (y_1, \dots, y_{i-1}, y_i, y_i + h, y_{i+1}, \dots, y_n).$$

Since $x_{i+1} = x_i + h$ and $0 \leq h \leq k$, and since, by induction hypothesis, y satisfies the statement, x also satisfies it.

Let us now show by induction on the length of the words that $\mathbf{FCat}^{(k)}$ contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, since $x_1 = 0$ and $0 \leq x_2 \leq k$, x has a factor $x_i x_{i+1}$ where i is the greatest integer such that $x_i \leq x_{i+1}$. Now, by setting $h := x_{i+1} - x_i$ and

$$(3.1.11) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n),$$

we have $x = y \circ_i 0h$. Since i is maximal, if $i + 2 \leq n$ we have $x_{i+2} < x_{i+1}$. This implies that y satisfies the statement. By induction hypothesis, $\mathbf{FCat}^{(k)}$ contains y and, since $x = y \circ_i 0h$, $\mathbf{FCat}^{(k)}$ also contains x . \square

A *k-leafy tree* is a planar rooted tree such that each internal node has exactly $k + 1$ children. The size $|T|$ of a *k-leafy tree* T is the number of its internal nodes. It is well-known that there are $\frac{1}{kn+1} \binom{kn+n}{n}$ *k-leafy trees* of size n . We say that an internal node x is *smaller* than an internal node y of T if, in the depth-first traversal of T , x appears before y . We also say that a *k-leafy tree* T is *well-labeled* if its root is labeled by 0, and, for each internal node x of T

labeled by \mathbf{a} , the children of x are labeled, from left to right, by $\mathbf{a} + k, \dots, \mathbf{a} + 1, \mathbf{a}$. There is a unique way to label a k -leafy tree so that it is well-labeled. There is a bijection $\phi_{\mathbf{FCat}^{(k)}}$ between the words of $\mathbf{FCat}^{(k)}$ of arity n and well-labeled k -leafy trees of size n .

To compute $\phi_{\mathbf{FCat}^{(k)}}(x)$ where x is an element of $\mathbf{FCat}^{(k)}$, iteratively insert the letters of x from left to right according to the following procedure. If $|x| = 1$, then $x = 0$ and $\phi_{\mathbf{FCat}^{(k)}}(x)$ is the only well-labeled k -leafy tree of size 1. Otherwise, the insertion of a letter $\mathbf{a} \geq 0$ into a well-labeled k -leafy tree T consists in replacing a leaf of T by the k -leafy tree S of size 1 labeled by \mathbf{a} so that S is the child of the greatest internal node such that the obtained k -tree is still well-labeled.

The inverse bijection is computed as follows. Given a well-labeled k -leafy tree T , one computes an element of $\mathbf{FCat}^{(k)}$ of arity n by reading its labels following a depth-first traversal of T .

Since the elements of $\mathbf{FCat}^{(k)}$ satisfy Proposition 3.7, $\phi_{\mathbf{FCat}^{(k)}}$ is well-defined. Hence, we can regard the elements of arity n of $\mathbf{FCat}^{(k)}$ as k -leafy trees of size n . Figure 4 shows an example of this bijection.

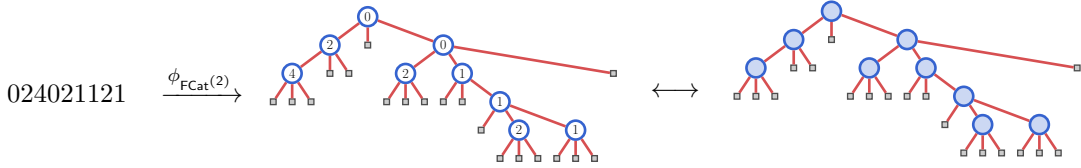


FIGURE 4. Interpretation of an element of the ns operad $\mathbf{FCat}^{(2)}$ in terms of 2-leafy trees via the bijection $\phi_{\mathbf{FCat}^{(2)}}$. The 2-leafy tree in the middle is well-labeled.

The bijection $\phi_{\mathbf{FCat}^{(k)}}$ between elements of $\mathbf{FCat}^{(k)}$ and k -leafy trees offers an alternative way to compute the composition of elements of $\mathbf{FCat}^{(k)}$:

Proposition 3.8. *Let S and T be two k -leafy trees and s be the i th internal node for the depth-first traversal of S . The composition $S \circ_i T$ in $\mathbf{FCat}^{(k)}$ amounts to replace s by the root of T and graft the children of s from right to left on the rightmost leaves of T .*

Proof. Let $x \in \mathbf{FCat}^{(k)}(n)$ and $y \in \mathbf{FCat}^{(k)}(m)$ such that $S := \phi_{\mathbf{FCat}^{(k)}}(x)$ and $T := \phi_{\mathbf{FCat}^{(k)}}(y)$. Let $U := \phi_{\mathbf{FCat}^{(k)}}(x \circ_i y)$. By definition of $\phi_{\mathbf{FCat}^{(k)}}$ and the partial composition maps of $\mathbf{FCat}^{(k)}$, U is obtained by inserting the prefix of length $i - 1$ of x , then the letters of y incremented by x_i , and finally, the suffix of length $n - i$ of x . Since by Proposition 3.7, y starts by 0, the internal nodes created by inserting the letters of y incremented by x_i are descendants of the internal node created by inserting $x_i = y_1$. Since the last $n - i$ letters of $x \circ_i y$ are the same as the last $n - i$ letters of x , by definition of $\phi_{\mathbf{FCat}^{(k)}}$, the children of the i th internal node of S are grafted in U from right to left on the rightmost leaves of T . This implies the statement. \square

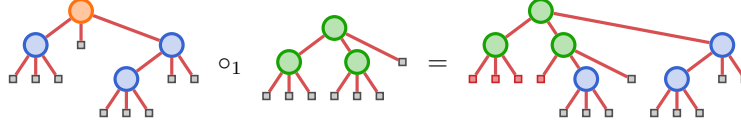


FIGURE 5. Interpretation of the partial composition map of the ns operad $\mathbf{FCat}^{(2)}$ in terms of 2-leafy trees.

Figure 5 shows an example of composition in $\mathbf{FCat}^{(2)}$.

The next Theorem elucidates the structure of $\mathbf{FCat}^{(k)}$:

Theorem 3.9. *The ns operad $\mathbf{FCat}^{(k)}$ admits the presentation*

$$(3.1.12) \quad \mathbf{FCat}^{(k)} = \mathcal{F}(\{\mathbf{a}_0, \dots, \mathbf{a}_k\}) / \equiv,$$

where the \mathbf{a}_i are of arity 2 and \equiv is the ns operadic congruence generated by

$$(3.1.13) \quad \mathbf{a}_{i+j} \circ_1 \mathbf{a}_i \leftrightarrow \mathbf{a}_i \circ_2 \mathbf{a}_j, \quad i, j \geq 0, i + j \leq k.$$

Proof. First, note that by replacing \mathbf{a}_i by $0i \in \mathbf{FCat}^{(k)}(2)$, we have $\text{ev}(x) = \text{ev}(y)$ for the relation $x \leftrightarrow y$ of the statement of the Theorem. Indeed, this equivalence class is the one of the element $(0, i, i + j)$ of $\mathbf{FCat}^{(k)}$.

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.14) \quad \begin{array}{c} \text{Diagram 1} \end{array} \mapsto \begin{array}{c} \text{Diagram 2} \end{array}.$$

This rewrite rule is terminating. Indeed, it is plain that for any rewriting $T_0 \rightarrow T_1$, we have $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all elements of $\mathcal{F}(\{\mathbf{a}_0, \dots, \mathbf{a}_k\})$ such that for each node x labeled by \mathbf{a}_i which has a left child y labeled by \mathbf{a}_j , one has $i < j$. This set \mathcal{S} of syntax trees admits the following regular specification

$$(3.1.15) \quad \mathcal{S} = \text{Diagram} + \bigsqcup_{0 \leq i \leq k} \mathcal{S}_i,$$

where \mathcal{S}_i is the set of such syntax trees with roots labeled by \mathbf{a}_i . These sets satisfy the following regular specification

$$(3.1.16) \quad \mathcal{S}_i = \begin{array}{c} \text{Diagram} \end{array} + \bigsqcup_{i+1 \leq j \leq k} \begin{array}{c} \text{Diagram} \end{array}.$$

Hence, the generating series $F(t)$ of \mathcal{S} and $F_i(t)$ of \mathcal{S}_i satisfy

$$(3.1.17) \quad F(t) = t + \sum_{0 \leq i \leq k} F_i(t),$$

and

$$(3.1.18) \quad F_i(t) = tF(t) + F(t) \sum_{i+1 \leq j \leq k} F_j(t).$$

By basic manipulations involving binomial coefficients, we obtain

$$(3.1.19) \quad F_i(t) = tF(t) \sum_{0 \leq j \leq k-i} \binom{k-i}{j} F(t)^j,$$

and then,

$$(3.1.20) \quad F(t) = t \sum_{0 \leq j \leq k+1} \binom{k+1}{j} F(t)^j.$$

The functional equation (3.1.20) is an alternative functional equation for the generating series of k -leafy trees. By Proposition 3.7, $F(t)$ also is the Hilbert series of $\mathbf{FCat}^{(k)}$.

Hence, by Lemma 1.2, $\mathbf{FCat}^{(k)}$ admits the claimed presentation. \square

3.1.4. *A ns operad on Schröder trees.* Let \mathbf{Schr} be the ns suboperad of \mathbf{TN} generated by 00, 01, and 10. The following table shows the first elements of \mathbf{Schr} .

Arity	Elements of \mathbf{Schr}
1	0
2	00, 01, 10
3	000, 001, 010, 011, 012, 021, 100, 101, 110, 120, 210
4	0000, 0001, 0010, 0011, 0012, 0021, 0100, 0101, 0110, 0111, 0112, 0120, 0121, 0122, 0123, 0132, 0210, 0211, 0212, 0221, 0231, 0321, 1000, 1001, 1010, 1011, 1012, 1021, 1100, 1101, 1110, 1120, 1200, 1201, 1210, 1220, 1230, 1320, 2100, 2101, 2110, 2120, 2210, 2310, 3210

Since $\mathbf{FCat}^{(1)}$ is generated by 00 and 01, $\mathbf{FCat}^{(1)}$ is a ns suboperad of \mathbf{Schr} . Moreover, since \mathbf{PW} is, by Proposition 3.2, generated as an operad by 00 and 01, \mathbf{Schr} is a ns suboperad of \mathbf{PW} .

One has the following characterization of the elements of \mathbf{Schr} :

Proposition 3.10. *The elements of \mathbf{Schr} are exactly the words x on the alphabet \mathbb{N} having at least one occurrence of 0 and, for all letter $\mathbf{b} \geq 1$ of x , there exists a letter $\mathbf{a} := \mathbf{b} - 1$ such that x has a factor \mathbf{aub} or \mathbf{bua} where u is a word consisting in letters \mathbf{c} satisfying $\mathbf{c} \geq \mathbf{b}$.*

Proof. Let us first show by induction on the length of the words that any word x of \mathbf{Schr} satisfies the statement. This is true when $|x| = 1$. When $|x| \geq 2$, by Lemma 1.1, there is a element y of \mathbf{Schr} of length $n := |x| - 1$, an integer $i \in [n]$, and $g \in \{00, 01, 10\}$ such that $x = y \circ_i g$. Then,

$$(3.1.21) \quad x = (y_1, \dots, y_{i-1}, y_i + g_1, y_i + g_2, y_{i+1}, \dots, y_{n-1}).$$

Since by induction hypothesis y satisfies the statement and $x_i = y_i$ or $x_{i+1} = y_i$, x also satisfies the statement.

Let us now show by induction on the length of the words that \mathbf{Schr} contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, let us observe that if x only consists in letters 0, \mathbf{Schr} contains x because x can be obtained by composing the generator 00 with itself. Hence, let us assume that x contains at least one letter different

from 0. Set \mathbf{b} as the greatest letter of x and $\mathbf{a} := \mathbf{b} - 1$. Since \mathbf{b} is the greatest letter of x , there is a factor $x_i x_{i+1}$ of x such that $x_i x_{i+1} \in \{\mathbf{bb}, \mathbf{ba}, \mathbf{ab}\}$. Set

$$(3.1.22) \quad y := (x_1, \dots, x_{i-1}, \min\{x_i, x_{i+1}\}, x_{i+2}, \dots, x_n),$$

and g as the generator 00 if $x_i = x_{i+1}$, as 01 if $x_i = x_{i+1} - 1$, or as 10 when $x_i = x_{i+1} + 1$. Then, we have $x = y \circ_i g$, and, since y is obtained from x by removing one of its greatest letter, y satisfies the statement. By induction hypothesis, \mathbf{Schr} contains y , and since $x = y \circ_i g$, \mathbf{Schr} also contains x . \square

A *Schröder tree* is a planar rooted tree such that no node has exactly one child. The size $|T|$ of a Schröder tree T is its number of leaves. There is a bijection $\phi_{\mathbf{Schr}}$ between the words of \mathbf{Schr} of arity n and Schröder trees of size n .

To compute $\phi_{\mathbf{Schr}}(x)$ where x is an element of \mathbf{Schr} , factorize x as $x = x^{(1)} \mathbf{a} \dots \mathbf{a} x^{(\ell)}$ where \mathbf{a} is the smallest letter occurring in x and the $x^{(i)}$ are factors of x without \mathbf{a} . Then, set

$$(3.1.23) \quad \phi_{\mathbf{Schr}}(x) := \begin{cases} \mathbf{1} & \text{if } x = \epsilon, \\ \bigwedge (\phi_{\mathbf{Schr}}(x^{(1)}), \dots, \phi_{\mathbf{Schr}}(x^{(\ell)})) & \text{otherwise,} \end{cases}$$

where ϵ denotes the empty word and $\bigwedge(T_1, \dots, T_\ell)$ is the Schröder tree consisting in a root that has T_1, \dots, T_ℓ as subtrees from left to right.

The inverse bijection is computed as follows. Given a Schröder tree T , one computes an element of \mathbf{Schr} by considering each internal node s and two adjacent consecutive edges of s and by assigning to these the depth of s . The element of \mathbf{Schr} is obtained by reading the labels from left to right.

Since the elements of \mathbf{Schr} satisfy Proposition 3.10, $\phi_{\mathbf{Schr}}$ is well-defined. Figure 6 shows an example of this bijection.

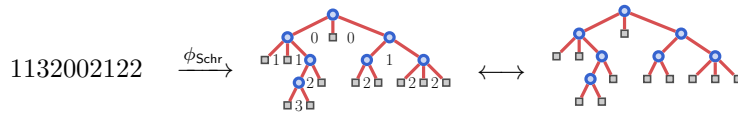


FIGURE 6. Interpretation of an element of the ns operad \mathbf{Schr} in terms of Schröder trees via the bijection $\phi_{\mathbf{Schr}}$.

Theorem 3.11. *The ns operad \mathbf{Schr} admits the presentation*

$$(3.1.24) \quad \mathbf{Schr} = \mathcal{F}(\{ \text{diagram 1}, \text{diagram 2}, \text{diagram 3} \}) / \equiv,$$

where $\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}$, $\begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix}$, and $\begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix}$ are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.1.25) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.26) \quad \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix},$$

$$(3.1.27) \quad \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.28) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.29) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.30) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix},$$

$$(3.1.31) \quad \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \circ_1 \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \leftrightarrow \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \circ_2 \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}.$$

Proof. First, note that by replacing $\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}$ by $00 \in \text{Schr}(2)$, $\begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix}$ by $01 \in \text{Schr}(2)$, and $\begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix}$ by $10 \in \text{Schr}(2)$, we have $\text{ev}(x) = \text{ev}(y)$ for the seven relations $x \leftrightarrow y$ of the statement of the Theorem. Indeed, then seven equivalence classes are, respectively, the ones of the elements 000, 101, 010, 001, 100, 011, and 110 of Schr .

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.32) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.33) \quad \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix},$$

$$(3.1.34) \quad \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.35) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.36) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix},$$

$$(3.1.37) \quad \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix},$$

$$(3.1.38) \quad \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \mapsto \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}.$$

This rewrite rule is terminating. Indeed, let T be an element of $\mathcal{F}(\{\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}, \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix}, \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix}\})$. By associating with T the pair $(k_T, w(T))$ where k_T is the number of nodes labeled by $\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}$ in T , it is plain that for any rewriting $T_0 \rightarrow T_1$, one has $k_{T_0} < k_{T_1}$ or $k_{T_0} = k_{T_1}$ and $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all elements of $\mathcal{F}(\{\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}, \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix}, \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix}\})$ such that nodes labeled by $\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}$ or $\begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix}$ have no left child and nodes labeled by $\begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix}$ have no right child labeled by $\begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix}$. This set \mathcal{S} of syntax trees admits the following regular specification

$$(3.1.39) \quad \mathcal{S} = \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} + \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \mathcal{S} + \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \mathcal{S} + \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagdown \\ \bullet \end{smallmatrix} \mathcal{S} + \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \mathcal{S} + \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \diagdown \\ \bullet \end{smallmatrix} \begin{smallmatrix} \bullet \\ \diagup \\ \bullet \end{smallmatrix} \mathcal{S}.$$

Hence, the generating series $F(t)$ of \mathcal{S} satisfies

$$(3.1.40) \quad F(t) = t + 3tF(t) + 2tF(t)^2,$$

that is the generating series of Schröder trees. By Proposition 3.10, $F(t)$ also is the Hilbert series of Schr .

Hence, by Lemma 1.2, Schr admits the claimed presentation. \square

3.1.5. *A ns operad on Motzkin words.* Let **Motz** be the ns suboperad of **TN** generated by 00 and 010. The following table shows the first elements of **Motz**.

Arity	Elements of Motz
1	0
2	00
3	000, 010
4	0000, 0010, 0100, 0110
5	00000, 00010, 00100, 00110, 01000, 01010, 01100, 01110, 01210
6	000000, 000010, 000100, 000110, 001000, 001010, 001100, 001110, 001210, 010000, 010010, 010100, 010110, 011000, 011010, 011100, 011110, 011210, 012100, 012110, 012210

Since 00 and 01 generate **FCat**⁽¹⁾ and since 010 = 00 ◦₁ 01, **Motz** is a ns suboperad of **FCat**⁽¹⁾. Moreover, since **FCat**⁽⁰⁾ is generated by 00, **FCat**⁽⁰⁾ is a ns suboperad of **Motz**.

One has the following characterization of the elements of **Motz**:

Proposition 3.12. *The elements of **Motz** are exactly the words x on the alphabet \mathbb{N} beginning and start by 0 and such that $|x_i - x_{i+1}| \leq 1$ for all $i \in [|x| - 1]$.*

Proof. Let us first show by induction on the length of the words that any word x of **Motz** satisfies the statement. This is true when $|x| = 1$. When $|x| \geq 2$, by Lemma 1.1, there is an element y of **Motz** of length $n < |x|$, an integer $i \in [n]$, and $g \in \{00, 010\}$ such that $x = y \circ_i g$. If $g = 00$, then one has

$$(3.1.41) \quad x = (y_1, \dots, y_{i-1}, y_i, y_i, y_{i+1}, \dots, y_n).$$

Otherwise, we have $g = 010$ and

$$(3.1.42) \quad x = (y_1, \dots, y_{i-1}, y_i, y_i + 1, y_i, y_{i+1}, \dots, y_n).$$

Since by induction hypothesis y satisfies the statement, it is immediate that in both cases, x also satisfies the statement.

Let us now show by induction on the length of the words that **Motz** contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, one has two cases to consider. If x contains a factor $x_i x_{i+1}$ such that $x_i = x_{i+1}$, by setting

$$(3.1.43) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n),$$

we have $x = y \circ_i 00$. Otherwise, let \mathbf{b} be the greatest letter of x . Since x satisfies the statement, there is in x a factor $x_i x_{i+1} x_{i+2}$ where $x_{i+1} = \mathbf{b}$ and $x_i = x_{i+2} = \mathbf{b} - 1$. By setting

$$(3.1.44) \quad y := (x_1, \dots, x_i, x_{i+3}, \dots, x_n),$$

we have $x = y \circ_i 010$. Now, for both cases, since by induction hypothesis, **Motz** contains y , **Motz** also contains x . \square

A *Motzkin word* is a word u on the alphabet $\{-1, 0, 1\}$ such that the sum of all letters of u is 0 and, for any prefix u' of u , the sum of all letters of u' is a nonnegative integer. The size $|u|$ of a Motzkin word u is its length plus one. In the sequel, we shall denote by $\bar{1}$ the letter -1 . We can represent a Motzkin word u graphically by a *Motzkin path* that is the path in \mathbb{N}^2 connecting the points $(0, 0)$ and $(n, 0)$ obtained by drawing a step $(1, -1)$ (resp. $(1, 0)$, $(1, 1)$) for each letter $\bar{1}$ (resp. 0 , 1) of u . There is a bijection ϕ_{Motz} between the words of **Motz** of arity n and Motzkin words of size n .

To compute $\phi_{\text{Motz}}(x)$ where x is an element of **Motz**(n), build the word u of length $n - 1$ satisfying $u_i := x_{i+1} - x_i$ for all $i \in [n - 1]$.

The inverse bijection is computed as follows. The element of **Motz** in bijection with a Motzkin word u is the word x such that x_i is the sum of the letters of the prefix $u_1 \dots u_{i-1}$ of u , for all $i \in [n]$.

Since the elements of **Motz** satisfy Proposition 3.12, ϕ_{Motz} is well-defined. Figure 7 shows an example of this bijection.

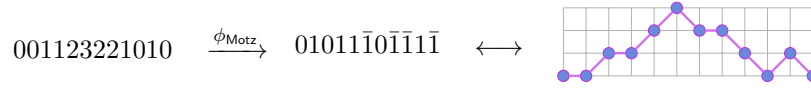


FIGURE 7. Interpretation of an element of the ns operad **Motz** in terms of Motzkin words and Motzkin paths via the bijection ϕ_{Motz} .

The bijection ϕ_{Motz} between elements of **Motz** and Motzkin words offers an alternative way to compute the composition of elements of **Motz**:

Proposition 3.13. *Let u and v be two Motzkin words where u is of size n , and $i \in [n]$ be an integer. Then the composition $u \circ_i v$ in **Motz** amounts to insert v at the i th position into u .*

Proof. Let y be the element of **Motz** in bijection by ϕ_{Motz} with v . The statement is a direct consequence of the fact that, by Proposition 3.12, y starts and ends by 0. \square

Figure 8 shows an example of composition in **Motz**.

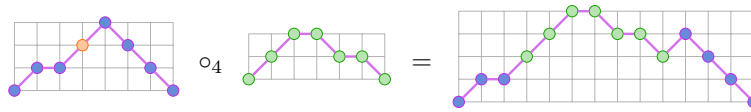


FIGURE 8. Interpretation of the partial composition map of the ns operad **Motz** in terms of Motzkin paths.

Theorem 3.14. *The ns operad \mathbf{Motz} admits the presentation*

$$(3.1.45) \quad \mathbf{Motz} = \mathcal{F}(\{\bullet\bullet, \bullet\bullet\bullet\}) / \equiv,$$

where $\bullet\bullet$ is of arity 2, $\bullet\bullet\bullet$ of arity 3, and \equiv is the ns operadic congruence generated by

$$(3.1.46) \quad \bullet\bullet \circ_1 \bullet\bullet \leftrightarrow \bullet\bullet \circ_2 \bullet\bullet, \quad (3.1.48) \quad \bullet\bullet \circ_1 \bullet\bullet\bullet \leftrightarrow \bullet\bullet\bullet \circ_3 \bullet\bullet,$$

$$(3.1.47) \quad \bullet\bullet\bullet \circ_1 \bullet\bullet \leftrightarrow \bullet\bullet \circ_2 \bullet\bullet\bullet, \quad (3.1.49) \quad \bullet\bullet\bullet \circ_1 \bullet\bullet\bullet \leftrightarrow \bullet\bullet\bullet \circ_3 \bullet\bullet\bullet.$$

Proof. First, note that by replacing $\bullet\bullet$ by $00 \in \mathbf{Motz}(2)$ and $\bullet\bullet\bullet$ by $010 \in \mathbf{Motz}(3)$, we have $\text{ev}(x) = \text{ev}(y)$ for the four relations $x \leftrightarrow y$ of the statement of the Theorem. Indeed, the four equivalence classes are, respectively, the ones of the elements 000, 0010, 0100, and 01010 of \mathbf{Motz} .

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.50) \quad \begin{array}{c} \text{Diagram 1} \end{array} \mapsto \begin{array}{c} \text{Diagram 2} \end{array}, \quad (3.1.52) \quad \begin{array}{c} \text{Diagram 3} \end{array} \mapsto \begin{array}{c} \text{Diagram 4} \end{array},$$

$$(3.1.51) \quad \begin{array}{c} \text{Diagram 5} \end{array} \mapsto \begin{array}{c} \text{Diagram 6} \end{array}, \quad (3.1.53) \quad \begin{array}{c} \text{Diagram 7} \end{array} \mapsto \begin{array}{c} \text{Diagram 8} \end{array}.$$

This rewrite rule is terminating. Indeed, it is plain that for any rewriting $T_0 \rightarrow T_1$, we have $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all syntax trees of $\mathcal{F}(\{\bullet\bullet, \bullet\bullet\bullet\})$ which have no internal node with an internal node as leftmost child. This set \mathcal{S} of trees admits the following regular specification

$$(3.1.54) \quad \mathcal{S} = \begin{array}{c} \text{Diagram 9} \end{array} + \begin{array}{c} \text{Diagram 10} \end{array} + \begin{array}{c} \text{Diagram 11} \end{array}.$$

Hence, the generating series $F(t)$ of \mathcal{S} satisfies

$$(3.1.55) \quad F(t) = t + tF(t) + tF(t)^2,$$

that is the generating series of Motzkin words. By Proposition 3.12, $F(t)$ also is the Hilbert series of \mathbf{Motz} .

Hence, by Lemma 1.2, \mathbf{Motz} admits the claimed presentation. \square

3.1.6. A ns operad on integer compositions. Let \mathbf{Comp} be the ns suboperad of \mathbf{TN}_2 generated by 00 and 01. The following table shows the first elements of \mathbf{Comp} .

Arity	Elements of \mathbf{Comp}
1	0
2	00, 01
3	000, 001, 010, 011
4	0000, 0001, 0010, 0011, 0100, 0101, 0110, 0111
5	00000, 00001, 00010, 00011, 00100, 00101, 00110, 00111, 01000, 01001, 01010, 01011, 01100, 01101, 01110, 01111

Since $\mathbf{FCat}^{(1)}$ is the ns suboperad of \mathbf{TN} generated by 00 and 01, and since \mathbf{TN}_2 is a quotient of \mathbf{TN} , \mathbf{Comp} is a quotient of $\mathbf{FCat}^{(1)}$. Moreover, $\mathbf{FCat}^{(0)}$ is a quotient of \mathbf{Comp} by the ns operadic congruence \equiv defined for all $x, y \in \mathbf{Comp}$ by $x \equiv y$.

One has the following characterization of the elements of \mathbf{Comp} :

Proposition 3.15. *The elements of \mathbf{Comp} are exactly the words on the alphabet $\{0, 1\}$ beginning by 0.*

Proof. It is immediate, from the definition of \mathbf{Comp} and Lemma 1.1, that any element of this ns operad starts by 0 since its generators 00 and 01 all start by 0.

Let us now show by induction on the length of the words that \mathbf{Comp} contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, let us observe that if x only consists in letters 0, \mathbf{Comp} contains x because x can be obtained by composing the generator 00 with itself. Otherwise, x has at least one occurrence of 1. Since its first letter is 0, there is in x a factor $x_i x_{i+1} = 01$. By setting

$$(3.1.56) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n),$$

we have $x = y \circ_i 01$. Since y satisfies the statement, by induction hypothesis \mathbf{Comp} contains y . Hence, \mathbf{Comp} also contains x . \square

An *integer composition* is a sequence $u_1 \dots u_k$ of positive integers. The size $|u|$ of an integer composition u is the sum of its letters. It is well-known that there are 2^{n-1} integer compositions of size n . We shall represent an integer composition $u := u_1 \dots u_k$ by a *ribbon diagram*, that is the diagram in which each letter u_i of u is encoded by a column consisting in u_i boxes, and the column encoding the letter u_{i+1} is attached on the right edge of the bottommost box of the column encoding u_i , for any $i \in [k-1]$. The *i th box* of a ribbon diagram D is the i th encountered box by traversing D column by column from left to right and from top to bottom. The *transpose* of D is the ribbon diagram obtained by applying on D the reflection through the line passing by its first and its last boxes. There is a bijection $\phi_{\mathbf{Comp}}$ between the words of \mathbf{Comp} of arity n and ribbon diagrams of integer compositions of size n .

To compute $\phi_{\mathbf{Comp}}(x)$ where x is an element of \mathbf{Comp} , iteratively insert the letters of x from left to right according to the following procedure. If $|x| = 1$, then $x = 0$ and $\phi_{\mathbf{Comp}}(0)$ is the only ribbon diagram consisting in one box. Otherwise, the insertion of a letter \mathbf{a} into D consists in adding a new box below (resp. to the right of) the right bottommost box of D if $\mathbf{a} = 1$ (resp. $\mathbf{a} = 0$).

The inverse bijection is computed as follows. Given a ribbon diagram D of an integer composition of size n , one computes an element of \mathbf{Comp} of arity n by labelling the first box of D by 0 and the i th box b by 0 if the $(i-1)$ st box is on the left of b or by 1 otherwise, for any $1 \leq i \leq n$. The corresponding element of \mathbf{Comp} is obtained by reading the labels of D from top to bottom and left to right.

Since the elements of \mathbf{Comp} satisfy Proposition 3.15, $\phi_{\mathbf{Comp}}$ is well-defined. Hence, we can regard the elements of arity n of \mathbf{Comp} as ribbon diagrams with n boxes. Figure 9 shows an example of this bijection.

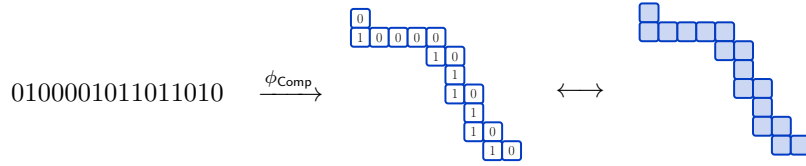


FIGURE 9. Interpretation of an element of the ns operad \mathbf{Comp} in terms of integer compositions via the bijection $\phi_{\mathbf{Comp}}$. Boxes of the ribbon diagram in the middle are labeled.

Encoding integer compositions by ribbon diagrams offers an alternative way to compute the composition of elements of \mathbf{Comp} :

Proposition 3.16. *Let C and D be two ribbon diagrams, i be an integer, and c be the i th box of C . Then, the composition $C \circ_i D$ in \mathbf{Comp} amounts to replace c by D if c is the upper box of its column, or to replace c by the transpose ribbon diagram of D otherwise.*

Proof. Let $x \in \mathbf{Comp}(n)$ and $y \in \mathbf{Comp}(m)$ such that $C := \phi_{\mathbf{Comp}}(x)$ and $D := \phi_{\mathbf{Comp}}(y)$. Let $E := \phi_{\mathbf{Comp}}(x \circ_i y)$. By definition of $\phi_{\mathbf{Comp}}$ and the partial composition maps of \mathbf{Comp} , E is obtained by inserting the prefix of length $i-1$ of x , then the letters of y incremented by x_i , and finally, the suffix of length $n-i$ of x . If c is the upper box of its column, then $x_i = 0$, and by definition of $\phi_{\mathbf{Comp}}$, E is obtained by replacing c by D in C . Otherwise, c is not the upper box of its column, and then $x_i = 1$. Immediately from the definition of $\phi_{\mathbf{Comp}}$, for any word z on the alphabet $\{0, 1\}$, the ribbon diagram $\phi_{\mathbf{Comp}}(0z)$ is the transpose of $\phi_{\mathbf{Comp}}(0\bar{z})$ where \bar{z} is the complementary of z . This implies the statement. \square

Figure 10 shows two examples of compositions in \mathbf{Comp} .

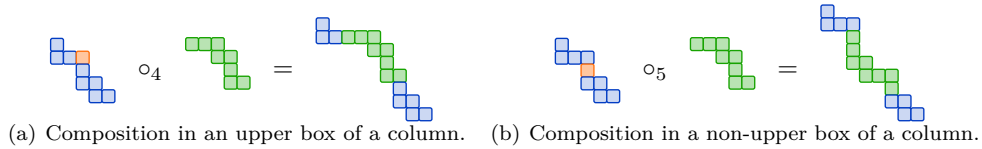


FIGURE 10. Interpretation of the partial composition map of the ns operad \mathbf{Comp} in terms of ribbon diagrams.

Theorem 3.17. *The ns operad \mathbf{Comp} admits the presentation*

$$(3.1.57) \quad \mathbf{Comp} = \mathcal{F}(\{\square, \blacksquare\}) / \equiv,$$

where \square and \blacksquare are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.1.58) \quad \square \circ_1 \square \leftrightarrow \square \circ_2 \square, \quad (3.1.60) \quad \blacksquare \circ_1 \blacksquare \leftrightarrow \blacksquare \circ_2 \square,$$

$$(3.1.59) \quad \blacksquare \circ_1 \square \leftrightarrow \square \circ_2 \blacksquare, \quad (3.1.61) \quad \square \circ_1 \blacksquare \leftrightarrow \blacksquare \circ_2 \blacksquare.$$

Proof. First, note that by replacing \square by $00 \in \mathbf{Comp}(2)$ and \blacksquare by $01 \in \mathbf{Comp}(2)$, we have $\text{ev}(x) = \text{ev}(y)$ for the four relations $x \leftrightarrow y$ of the statement of the Theorem. Indeed, the four equivalence classes are, respectively, the ones of the elements 000, 001, 011, and 100 of \mathbf{Comp} .

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.62) \quad \begin{array}{c} \square \\ \swarrow \searrow \\ \square \quad \square \end{array} \mapsto \begin{array}{c} \square \\ \swarrow \searrow \\ \square \quad \square \end{array}, \quad (3.1.64) \quad \begin{array}{c} \blacksquare \\ \swarrow \searrow \\ \square \quad \square \end{array} \mapsto \begin{array}{c} \blacksquare \\ \swarrow \searrow \\ \square \quad \square \end{array},$$

$$(3.1.63) \quad \begin{array}{c} \blacksquare \\ \swarrow \searrow \\ \square \quad \square \end{array} \mapsto \begin{array}{c} \blacksquare \\ \swarrow \searrow \\ \square \quad \square \end{array}, \quad (3.1.65) \quad \begin{array}{c} \square \\ \swarrow \searrow \\ \square \quad \square \end{array} \mapsto \begin{array}{c} \square \\ \swarrow \searrow \\ \square \quad \square \end{array}.$$

This rewrite rule is terminating. Indeed, it is plain that for any rewriting $T_0 \rightarrow T_1$, we have $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all syntax trees of $\mathcal{F}(\{\square, \blacksquare\})$ which have no internal node with an internal node as leftmost child. Hence, the generating series $F(t)$ of the normal forms of \mapsto is

$$(3.1.66) \quad F(t) = \sum_{n \geq 1} 2^{n-1} t^n.$$

By Proposition 3.15, $F(t)$ also is the Hilbert series of \mathbf{Comp} .

Hence, by Lemma 1.2, \mathbf{Comp} admits the claimed presentation. \square

3.1.7. A ns operad on directed animals. Let \mathbf{DA} be the ns suboperad of \mathbf{TN}_3 generated by 00 and 01. We shall here denote by $\bar{1}$ the representative of the equivalence class of 2 in \mathbb{N}_3 . The following table shows the first elements of \mathbf{DA} .

Arity	Elements of \mathbf{DA}
1	0
2	00, 01
3	000, 001, 010, 011, 01 $\bar{1}$
4	0000, 0001, 0010, 0011, 001 $\bar{1}$, 0100, 0101, 0110, 0111, 011 $\bar{1}$, 01 $\bar{1}$ 0, 01 $\bar{1}$ 1, 01 $\bar{1}\bar{1}$
5	00000, 00001, 00010, 00011, 0001 $\bar{1}$, 00100, 00101, 00110, 00111, 0011 $\bar{1}$, 001 $\bar{1}$ 0, 001 $\bar{1}$ 1, 001 $\bar{1}\bar{1}$, 01000, 01001, 01010, 01011, 0101 $\bar{1}$, 01100, 01101, 01110, 01111, 0111 $\bar{1}$, 011 $\bar{1}$ 0, 011 $\bar{1}$ 1, 011 $\bar{1}\bar{1}$, 01 $\bar{1}$ 00, 01 $\bar{1}$ 01, 01 $\bar{1}$ 0 $\bar{1}$, 01 $\bar{1}$ 10, 01 $\bar{1}$ 11, 01 $\bar{1}$ 1 $\bar{1}$, 01 $\bar{1}\bar{1}$ 0, 01 $\bar{1}\bar{1}$ 1, 01 $\bar{1}\bar{1}\bar{1}$

Since $\mathbf{FCat}^{(1)}$ is the ns suboperad of \mathbf{TN} generated by 00 and 01, and since \mathbf{TN}_3 is a quotient of \mathbf{TN} , \mathbf{DA} is a quotient of $\mathbf{FCat}^{(1)}$. Moreover, \mathbf{DA} is a quotient of $\mathbf{FCat}^{(0)}$ by the ns operadic congruence \equiv defined for all $x, y \in \mathbf{DA}$ by $x \equiv y$.

Proposition 3.18. *Let $\phi_{\mathbf{DA}} : \mathbf{DA}(n) \rightarrow \{\bar{1}, 0, 1\}^{n-1}$ the mapping defined for any element x of arity n of \mathbf{DA} by*

$$(3.1.67) \quad \phi_{\mathbf{DA}}(x) := (x_1 * x_2, x_2 * x_3, \dots, x_{n-1} * x_n),$$

where $x_i * x_{i+1} := x_{i+1} - x_i \pmod{3}$. Then, $\phi_{\mathbf{DA}}$ is a bijection between the elements of arity n of \mathbf{DA} and prefixes of Motzkin words of length $n - 1$.

Proof. Let us first show by induction on the length of the words that for any $x \in \mathbf{DA}$, $\phi_{\mathbf{DA}}(x)$ is a prefix of a Motzkin word of length $|x| - 1$. This is true when $|x| = 1$. When $|x| \geq 2$, by Lemma 1.1, there is an element y of \mathbf{DA} of length $n := |x| - 1$, an integer $i \in [n]$, and $g \in \{00, 01\}$ such that $x = y \circ_i g$. We now have two cases depending on g .

Case 1. If $g = 00$, then

$$(3.1.68) \quad x = (y_1, \dots, y_i, y_i, y_{i+1}, \dots, y_n).$$

By induction hypothesis, $\phi_{\mathbf{DA}}(y)$ is a prefix a Motzkin word of length $n - 1$. Since x is obtained from y by duplicating its i th letter, $\phi_{\mathbf{DA}}(x)$ is obtained from $\phi_{\mathbf{DA}}(y)$ by inserting a 0 at an appropriate place. Hence, $\phi_{\mathbf{DA}}(x)$ is a prefix of a Motzkin word of length n .

Case 2. Otherwise, we have $g = 01$ and then,

$$(3.1.69) \quad x = (y_1, \dots, y_i, y_i + 1, y_{i+1}, \dots, y_n),$$

where $+$ denotes the addition in \mathbb{N}_3 . We have now two sub-cases whether y_i is the last letter of y .

Case 2.1. If it is the case, then $x_{i+1} = y_i + 1$ is the last letter of x and $\phi_{\mathbf{DA}}(x)$ is obtained from $\phi_{\mathbf{DA}}(y)$ by concatenating a 1 on the right. Hence, since by induction hypothesis, $\phi_{\mathbf{DA}}(y)$ is a prefix of a Motzkin word of length $n - 1$, $\phi_{\mathbf{DA}}(x)$ is a prefix of a Motzkin word of length n .

Case 2.2. Otherwise, we have $i < n$. We observe that $\phi_{\mathbf{DA}}(x)$ is obtained from $\phi_{\mathbf{DA}}(y)$ by replacing a letter 0 (resp. 1, $\bar{1}$) by a factor $1\bar{1}$ (resp. 10, 11) at an appropriate place. Hence, since by induction hypothesis, $\phi_{\mathbf{DA}}(y)$ is a prefix of a Motzkin word of length $n - 1$, $\phi_{\mathbf{DA}}(x)$ is a prefix of a Motzkin word of length n .

Let us now show that $\phi_{\mathbf{DA}}$ is a bijection between the elements of arity n of \mathbf{DA} and prefixes of Motzkin words of length $n - 1$.

The injectivity of $\phi_{\mathbf{DA}}$ is a direct consequence of the fact that, given a element x of \mathbf{DA} and a letter $\mathbf{a} \in \{\bar{1}, 0, 1\}$, there is at most one letter $\mathbf{b} \in \{\bar{1}, 0, 1\}$ such that $\phi_{\mathbf{DA}}(x\mathbf{b}) = \phi_{\mathbf{DA}}(x)\mathbf{a}$.

Let us finally show that $\phi_{\mathbf{DA}}$ is a surjection. We proceed by induction on the length of the words to construct for any prefix of a Motzkin word u an element x of \mathbf{DA} such that $\phi_{\mathbf{DA}}(x) = u$. When u is the empty word, $x := 0$ is an element of $\mathbf{DA}(1)$ and since $\phi_{\mathbf{DA}}(x)$ is the empty word, the property is satisfied. When $n := |u| \geq 1$, one has two cases to consider depending on the last letter u_n of u .

Case 1'. If $u_n \in \{0, 1\}$, by induction hypothesis, there is an element y of $\text{DA}(n)$ such that $\phi_{\text{DA}}(y) = u_1 \dots u_{n-1}$. Hence, by setting $x := y \circ_n 0u_n$, it follows, by definition of ϕ_{DA} , that x is a preimage of u for ϕ_{DA} .

Case 2'. Otherwise, we have $u_n = \bar{1}$ and there is at least one occurrence of a 1 in u . Hence, let $i \in [n-1]$ be the greatest integer such that $u_i = 1$. We now have two sub-cases depending on the value of u_{i+1} .

Case 2'.1. If $u_{i+1} = 0$, the word

$$(3.1.70) \quad u' := u_1 \dots u_i u_{i+2} \dots u_n$$

is still a prefix of a Motzkin word. Then, by induction hypothesis, there is an element y of $\text{DA}(n)$ such that $\phi_{\text{DA}}(y) = u'$. Hence, by setting $x := y \circ_i 00$, it follows, by definition of ϕ_{DA} , that x is a preimage of u for ϕ_{DA} .

Case 2'.2. Otherwise, we have $u_{i+1} = \bar{1}$. Then, the word

$$(3.1.71) \quad u' := u_1 \dots u_{i-1} 0u_{i+2} \dots u_n$$

is still a prefix of a Motzkin word. Then, by induction hypothesis, there is an element y of $\text{DA}(n)$ such that $\phi_{\text{DA}}(y) = u'$. Hence by setting $x := y \circ_i 01$, it follows, by definition of ϕ_{DA} , that x is a preimage of u for ϕ_{DA} .

We then have proved that ϕ_{DA} is well-defined, injective, and surjective. Hence, it is a bijection between elements of arity n of DA and prefixes of Motzkin words of length $n-1$. \square

Here are two examples of images by ϕ_{DA} of elements of DA .

$$(3.1.72) \quad \phi_{\text{DA}}(011\bar{1}\bar{1}0\bar{1}01) = 10101\bar{1}\bar{1}1,$$

$$(3.1.73) \quad \phi_{\text{DA}}(010010101\bar{1}1) = 1\bar{1}01\bar{1}\bar{1}11\bar{1}.$$

Recall that a *directed animal* is a subset A of \mathbb{N}^2 such that $(0, 0) \in A$ and $(i, j) \in A$ with $i \geq 1$ or $j \geq 1$ implies $(i-1, j) \in A$ or $(i, j-1) \in A$. The size of a directed animal A is its cardinality. Figure 11 shows a directed animal.

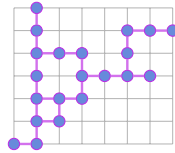


FIGURE 11. A directed animal of size 21. The point $(0, 0)$ is the lowest and leftmost point.

According to [GBV88], there is a bijection α between the set of prefixes of Motzkin words of length $n-1$ and the set of directed animals of size n . Hence, by Proposition 3.18, the map $\alpha \circ \phi_{\text{DA}}$ is a bijection between the elements of DA of arity n and directed animals of size n and moreover, DA can be seen as a ns operad on directed animals.

Theorem 3.19. *The ns operad DA admits the presentation*

$$\text{DA} = \mathcal{F}(\{\bullet\bullet, \bullet\bullet\}) / \equiv,$$

where $\bullet\bullet$ and $\bullet\bullet$ are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.1.74) \quad \bullet\bullet \circ_1 \bullet\bullet \leftrightarrow \bullet\bullet \circ_2 \bullet\bullet, \quad (3.1.76) \quad \bullet\bullet \circ_1 \bullet\bullet \leftrightarrow \bullet\bullet \circ_2 \bullet\bullet,$$

$$(3.1.75) \quad \bullet\bullet \circ_1 \bullet\bullet \leftrightarrow \bullet\bullet \circ_2 \bullet\bullet, \quad (3.1.77) \quad (\bullet\bullet \circ_1 \bullet\bullet) \circ_2 \bullet\bullet \leftrightarrow (\bullet\bullet \circ_2 \bullet\bullet) \circ_3 \bullet\bullet.$$

Proof. First, note that by replacing $\bullet\bullet$ by $00 \in \text{DA}(2)$ and $\bullet\bullet$ by $01 \in \text{DA}(2)$, we have $\text{ev}(x) = \text{ev}(y)$ for the four relations $x \leftrightarrow y$ of the statement of the Theorem. Indeed, the four equivalence classes are, respectively, the ones of the elements 000 , 001 , 011 , and $01\bar{1}0$ of DA .

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.78) \quad \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \bullet\bullet \quad \bullet\bullet \end{array} \mapsto \begin{array}{c} \bullet\bullet \\ \diagdown \quad \diagup \\ \bullet\bullet \quad \bullet\bullet \end{array}, \quad (3.1.80) \quad \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \bullet\bullet \quad \bullet\bullet \end{array} \mapsto \begin{array}{c} \bullet\bullet \\ \diagdown \quad \diagup \\ \bullet\bullet \quad \bullet\bullet \end{array},$$

$$(3.1.79) \quad \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \bullet\bullet \quad \bullet\bullet \end{array} \mapsto \begin{array}{c} \bullet\bullet \\ \diagdown \quad \diagup \\ \bullet\bullet \quad \bullet\bullet \end{array}, \quad (3.1.81) \quad \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \bullet\bullet \quad \bullet\bullet \end{array} \mapsto \begin{array}{c} \bullet\bullet \\ \diagdown \quad \diagup \\ \bullet\bullet \quad \bullet\bullet \end{array}.$$

This rewrite rule is terminating. Indeed, let T be a syntax tree of $\mathcal{F}(\{\bullet\bullet, \bullet\bullet\})$. By associating the pair $(-k_T, w(T))$ with T , where k_T is the sum, for all internal nodes x of T labeled by $\bullet\bullet$, of the number of internal nodes constituting the right subtree of x , it is plain that for any rewriting $T_0 \rightarrow T_1$, we have $k_{T_0} < k_{T_1}$, or $k_{T_0} = k_{T_1}$ and $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all syntax trees of $\mathcal{F}(\{\bullet\bullet, \bullet\bullet\})$ such that no internal node labeled by $\bullet\bullet$ has a left child labeled by $\bullet\bullet$, no internal node labeled by $\bullet\bullet$ has a child labeled by $\bullet\bullet$, and no internal node labeled by $\bullet\bullet$ has a right child labeled by $\bullet\bullet$ which has a right child labeled by $\bullet\bullet$. This set \mathcal{S} of syntax trees admits the following regular specification

$$(3.1.82) \quad \mathcal{S} = \mathcal{T} + \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \tau \quad s \end{array},$$

where \mathcal{T} is the set of syntax trees admitting the following regular specification

$$(3.1.83) \quad \mathcal{T} = \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \tau \quad \tau \end{array} + \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \tau \quad \tau \end{array} + \begin{array}{c} \bullet\bullet \\ \diagup \quad \diagdown \\ \tau \quad \tau \end{array}.$$

Hence, the generating series $F(t)$ of \mathcal{S} satisfies

$$(3.1.84) \quad F(t) = \frac{1 - 3t - (1 - 2t - 3t^2)^{1/2}}{6t - 2},$$

which is the generating function of directed animals. By Proposition 3.18, $F(t)$ also is the Hilbert series of DA .

Hence, by Lemma 1.2, DA admits the claimed presentation. \square

Since the nontrivial relation (3.1.77) has degree 3, the presentation of DA exhibited by Theorem 3.19 is not quadratic. Moreover, DA is not a quadratic ns operad since, as an

exhaustive inspection can show, there is no quadratic ns operad generated by two generators of arity 2 which has the same dimensions as DA.

3.1.8. *A ns operad on segmented integer compositions.* Let **SComp** be the ns suboperad of \mathbf{TN}_3 generated by 00, 01, and 02. The following table shows the first elements of **SComp**.

Arity	Elements of SComp
1	0
2	00, 01, 02
3	000, 001, 002, 010, 011, 012, 020, 021, 022
4	0000, 0001, 0002, 0010, 0011, 0012, 0020, 0021, 0022, 0100, 0101, 0102, 0110, 0111, 0112, 0120, 0121, 0122, 0200, 0201, 0202, 0210, 0211, 0212, 0220, 0221, 0222

Since $\mathbf{FCat}^{(2)}$ is the ns suboperad of \mathbf{TN} generated by 00, 01, and 02, and since \mathbf{TN}_3 is a quotient of \mathbf{TN} , **SComp** is a quotient of $\mathbf{FCat}^{(2)}$. Moreover, since DA is generated by 00 and 01, DA is a ns suboperad of **SComp**.

One has the following characterization of the elements of **SComp**:

Proposition 3.20. *The elements of **SComp** are exactly the words on the alphabet $\{0, 1, 2\}$ beginning by 0.*

Proof. It is immediate, from the definition of **SComp** and Lemma 1.1, that any element of this ns operad starts by 0 since its generators 00, 01, and 02 all start by 0.

Let us now show by induction on the length of the words that **SComp** contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, let us observe that if x only consists in letters 0, **SComp** contains x because x can be obtained by composing the generator 00 with itself. Otherwise, x has at least one occurrence of a 1 or a 2. Since its first letter is 0, there is in x a factor $x_i x_{i+1} =: g$ such that $g \in \{01, 02\}$. By setting

$$(3.1.85) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n),$$

we have $x = y \circ_i g$. Since y satisfies the statement, by induction hypothesis **SComp** contains y . Hence, **SComp** also contains x . \square

A *segmented integer composition* is a sequence (S_1, \dots, S_ℓ) of integer compositions. The size $|S|$ of a segmented integer composition is the sum of the sizes of the integer compositions which constitute S . It is well-known that there are 3^{n-1} segmented integer compositions of size n . We shall represent a segmented integer composition S by a *ribbon diagram*, that is the diagram consisting in the sequence of the ribbon diagrams of the integer compositions that constitute S . There is a bijection between the words of **SComp** of arity n and ribbon diagrams of segmented compositions of size n .

To compute $\phi_{\mathbf{SComp}}(x)$ where x is an element of **SComp**, factorize x as $x = 0x^{(1)} \dots 0x^{(\ell)}$ such that for any $i \in [\ell]$, the factor $x^{(i)}$ has no occurrence of 0, and compute the sequence

$(\phi_{\mathbf{Comp}}(0\bar{x}^{(1)}), \dots, \phi_{\mathbf{Comp}}(0\bar{x}^{(\ell)}))$, where for any $i \in [\ell]$, $\bar{x}^{(i)}$ is the word obtained from $x^{(i)}$ by decreasing all letters.

The inverse bijection is computed as follows. Given a ribbon diagram $S := (S_1, \dots, S_\ell)$ of a segmented integer composition of size n , one computes an element of \mathbf{SComp} of arity n by computing the sequence $(u^{(1)}, \dots, u^{(\ell)})$ where for any $i \in [\ell]$, $u^{(i)}$ is the word of \mathbf{Comp} obtained by applying the inverse bijection of $\phi_{\mathbf{Comp}}$ on $u^{(i)}$, then by incrementing in each $u^{(i)}$ all letters, excepted the first one, and finally by concatenating each words of the sequence.

Since the elements of \mathbf{SComp} satisfy Proposition 3.20, $\phi_{\mathbf{SComp}}$ is well-defined. Figure 12 shows an example of this bijection.

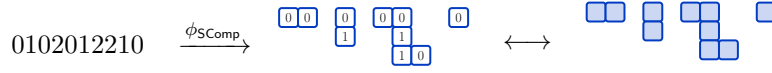


FIGURE 12. Interpretation of an element of the operad \mathbf{SComp} in terms of a segmented composition via the bijection $\phi_{\mathbf{SComp}}$. Boxes of the ribbon diagram in the middle are labeled.

Theorem 3.21. *The ns operad \mathbf{SComp} admits the presentation*

$$(3.1.86) \quad \mathbf{SComp} = \mathcal{F}(\{\text{red square}, \text{blue square}, \text{orange square}\}) / \equiv,$$

where red square , blue square , and orange square are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.1.87) \quad \text{red square} \circ_1 \text{red square} \leftrightarrow \text{red square} \circ_2 \text{red square},$$

$$(3.1.88) \quad \text{blue square} \circ_1 \text{red square} \leftrightarrow \text{red square} \circ_2 \text{blue square},$$

$$(3.1.89) \quad \text{blue square} \circ_1 \text{blue square} \leftrightarrow \text{blue square} \circ_2 \text{red square},$$

$$(3.1.90) \quad \text{red square} \circ_1 \text{blue square} \leftrightarrow \text{blue square} \circ_2 \text{orange square},$$

$$(3.1.91) \quad \text{blue square} \circ_1 \text{orange square} \leftrightarrow \text{orange square} \circ_2 \text{orange square},$$

$$(3.1.92) \quad \text{red square} \circ_1 \text{orange square} \leftrightarrow \text{orange square} \circ_2 \text{blue square},$$

$$(3.1.93) \quad \text{orange square} \circ_1 \text{red square} \leftrightarrow \text{red square} \circ_2 \text{orange square},$$

$$(3.1.94) \quad \text{orange square} \circ_1 \text{blue square} \leftrightarrow \text{blue square} \circ_2 \text{blue square},$$

$$(3.1.95) \quad \text{orange square} \circ_1 \text{orange square} \leftrightarrow \text{orange square} \circ_2 \text{red square}.$$

Proof. First, note that by replacing red square by $00 \in \mathbf{SComp}(2)$, blue square by $01 \in \mathbf{SComp}(2)$, and orange square by $02 \in \mathbf{SComp}(2)$, we have $\text{ev}(x) = \text{ev}(y)$ for the nine relations $x \leftrightarrow y$ of the statement of the Theorem. Indeed, the nine equivalence classes are, respectively, the ones of the elements 000, 001, 010, 021, 020, 002, 012, and 022.

Consider now the orientation of \leftrightarrow into the rewrite rule \mapsto defined by

$$(3.1.96) \quad \begin{array}{c} \text{Diagram 1} \end{array} \rightarrow \begin{array}{c} \text{Diagram 2} \end{array},$$

$$(3.1.97) \quad \begin{array}{c} \text{Diagram 3} \end{array} \rightarrow \begin{array}{c} \text{Diagram 4} \end{array},$$

$$(3.1.98) \quad \begin{array}{c} \text{Diagram 5} \end{array} \rightarrow \begin{array}{c} \text{Diagram 6} \end{array},$$

$$(3.1.99) \quad \begin{array}{c} \text{Diagram 7} \end{array} \rightarrow \begin{array}{c} \text{Diagram 8} \end{array},$$

$$(3.1.100) \quad \begin{array}{c} \text{Diagram 9} \end{array} \rightarrow \begin{array}{c} \text{Diagram 10} \end{array},$$

$$(3.1.101) \quad \begin{array}{c} \text{Diagram 11} \end{array} \rightarrow \begin{array}{c} \text{Diagram 12} \end{array},$$

$$(3.1.102) \quad \begin{array}{c} \text{Diagram 13} \end{array} \rightarrow \begin{array}{c} \text{Diagram 14} \end{array},$$

$$(3.1.103) \quad \begin{array}{c} \text{Diagram 15} \end{array} \rightarrow \begin{array}{c} \text{Diagram 16} \end{array},$$

$$(3.1.104) \quad \begin{array}{c} \text{Diagram 17} \end{array} \rightarrow \begin{array}{c} \text{Diagram 18} \end{array}.$$

This rewrite rule is terminating. Indeed, it is plain that for any rewriting $T_0 \rightarrow T_1$, we have $w(T_0) < w(T_1)$.

Moreover, the normal forms of \mapsto are all syntax trees of $\mathcal{F}(\{\text{red square}, \text{blue square}, \text{orange square}\})$ such that each internal node has no internal node as left son. Hence, the generating series $F(t)$ of the normal forms of \mapsto is

$$(3.1.105) \quad F(t) = \sum_{n \geq 1} 3^{n-1} t^n.$$

By Proposition 3.20, $F(t)$ also is the Hilbert series of **SComp**.

Hence, by Lemma 1.2, **SComp** admits the claimed presentation. \square

3.2. Operads from the multiplicative monoid. We shall denote by **M** the multiplicative monoid of integers.

Note that the ns suboperad of **TM** generated by 00 and the ns suboperad of **TM** generated by 11 are both isomorphic to the associative commutative operad **Com**.

The operads constructed in this section fit into the diagram of ns operads represented by Figure 13. Table 2 summarizes some information about these ns operads.



FIGURE 13. The diagram of ns suboperads and quotients of **TM**. Arrows \succrightarrow (resp. \twoheadrightarrow) are injective (resp. surjective) ns operad morphisms.

Monoid	Ns operad	Generators	First dimensions	Combinatorial objects
\mathbb{M}	Tr	01, 10, 11	1, 3, 7, 15, 31	Binary words with at least one 1
	Di	01, 10	1, 2, 3, 4, 5	Binary words with exactly one 1

TABLE 2. Ground monoids, generators, first dimensions, and combinatorial objects involved in the ns suboperads and quotients of $\mathbb{T}\mathbb{M}$.

3.2.1. *The diassociative operad.* Let Di be the ns suboperad of $\mathbb{T}\mathbb{M}$ generated by 01 and 10. The following table shows the first elements of Di .

Arity	Elements of Di
1	1
2	01, 10
3	001, 010, 100
4	0001, 0010, 0100, 1000
5	00001, 00010, 00100, 01000, 10000
6	000001, 000010, 000100, 001000, 010000, 100000

Note that Com is a quotient of Di by the ns operadic congruence \equiv defined for all $x, y \in \text{Di}(n)$ by $x \equiv y$.

One has the following characterization of the elements of Di :

Proposition 3.22. *The elements of Di are exactly the words on the alphabet $\{0, 1\}$ containing exactly one 1.*

Proof. Let us first show by induction on the length of the words that any word x of Di satisfies the statement. This is true when $|x| = 1$ since 1 is the unit of \mathbb{M} . When $|x| \geq 2$, by Lemma 1.1, there is an element y of Di of length $n := |x| - 1$, an integer $i \in [n]$, and $g \in \{01, 10\}$ such that $x = y \circ_i g$. In all cases, x is obtained from y by inserting a 0 at an appropriate position. Since, by induction hypothesis, y satisfies the statement, x also satisfies the statement.

Let us now show by induction on the length of the words that Di contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, there is in x a factor $x_i x_{i+1} =: g$ such that $g \in \{01, 10\}$. Assume without loss of generality that $g = 01$. Then, by setting

$$(3.2.1) \quad y := (x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n),$$

we have $x = y \circ_i g$. Since y satisfies the statement, by induction hypothesis Di contains y . Hence, Di also contains x . \square

Recall that the *diassociative operad* [Lod01] Dias is the ns operad admitting the presentation

$$(3.2.2) \quad \text{Dias} := \mathcal{F}(\{\neg, \vdash\}) / \equiv,$$

where \neg and \vdash are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.2.3) \quad \dashv \circ_1 \dashv \leftrightarrow \dashv \circ_2 \dashv \leftrightarrow \dashv \circ_2 \vdash,$$

$$(3.2.5) \quad \dashv \circ_1 \vdash \leftrightarrow \vdash \circ_2 \dashv.$$

$$(3.2.4) \quad \vdash \circ_2 \vdash \leftrightarrow \vdash \circ_1 \vdash \leftrightarrow \vdash \circ_1 \dashv,$$

Proposition 3.23. *The ns operads \mathbf{Di} and \mathbf{Dias} are isomorphic and the map*

$$(3.2.6) \quad \phi : \mathbf{Dias} \rightarrow \mathbf{Di}$$

satisfying $\phi(\dashv) = 10$ and $\phi(\vdash) = 01$ is an isomorphism.

Proof. By replacing each generator g of \mathbf{Dias} by $\phi(g)$, the generators 01 and 10 of \mathbf{Di} satisfy the same relations as the generators \dashv and \vdash of \mathbf{Dias} . Moreover, by Proposition 3.22, the Hilbert series of \mathbf{Di} is

$$(3.2.7) \quad F(t) = \sum_{n \geq 1} nt^n,$$

which also is the Hilbert series of \mathbf{Dias} . Then, there is no nontrivial relation of degree greater than two involving generators of \mathbf{Di} . \square

Proposition 3.23 also shows that \mathbf{Di} is a realization of the diassociative operad.

3.2.2. The triassociative operad. Let \mathbf{Tr} be the ns suboperad of $\mathbf{T\mathbb{M}}$ generated by 01 , 10 , and 11 . The following table shows the first elements of \mathbf{Tr} .

Arity	Elements of \mathbf{Tr}
1	1
2	01, 10, 11
3	001, 010, 011, 100, 101, 110, 111
4	0001, 0010, 0011, 0100, 0101, 0110, 0111, 1000, 1001, 1010, 1011, 1100, 1101, 1110, 1111

Since \mathbf{Di} is generated by 01 and 10 , \mathbf{Di} is a ns suboperad of \mathbf{Tr} .

One has the following characterization of the elements of \mathbf{Tr} :

Proposition 3.24. *The elements of \mathbf{Tr} are exactly the words on the alphabet $\{0, 1\}$ containing at least one 1.*

Proof. Let us first show by induction on the length of the words that any word x of \mathbf{Tr} satisfies the statement. This is true when $|x| = 1$ since 1 is the unit of \mathbb{M} . When $|x| \geq 2$, by Lemma 1.1, there is an element y of \mathbf{Tr} of length $n := |x| - 1$, an integer $i \in [n]$, and $g \in \{01, 10, 11\}$ such that $x = y \circ_i g$. By induction hypothesis, y contains at least one 1. Since all generators of \mathbf{Tr} contain at least one 1, x also contains at least one 1.

Let us now show by induction on the length of the words that \mathbf{Tr} contains any word x satisfying the statement. This is true when $|x| = 1$. When $n := |x| \geq 2$, there is in x a factor $x_i x_{i+1} =: g$ such that $g \in \{01, 10, 11\}$. Then, by setting

$$(3.2.8) \quad y := (x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n)$$

if $g = 01$, or

$$(3.2.9) \quad y := (x_1, \dots, x_i, x_{i+2}, \dots, x_n)$$

if $g \in \{10, 11\}$, we have $x = y \circ_i g$. Since y satisfies the statement, by induction hypothesis \mathbf{Tr} contains y . Hence, \mathbf{Tr} also contains x . \square

Recall that the *triassociative operad* [LR04] \mathbf{Trias} is the ns operad admitting the presentation

$$(3.2.10) \quad \mathbf{Trias} := \mathcal{F}(\{\neg, \perp, \vdash\}) / \equiv,$$

where \neg , \perp , and \vdash are of arity 2, and \equiv is the ns operadic congruence generated by

$$(3.2.11) \quad \neg \circ_1 \vdash \leftrightarrow \vdash \circ_2 \neg,$$

$$(3.2.15) \quad \perp \circ_1 \vdash \leftrightarrow \vdash \circ_2 \perp,$$

$$(3.2.12) \quad \perp \circ_1 \perp \leftrightarrow \perp \circ_2 \perp,$$

$$(3.2.16) \quad \neg \circ_1 \neg \leftrightarrow \neg \circ_2 \neg \leftrightarrow \neg \circ_2 \vdash \leftrightarrow \neg \circ_2 \perp,$$

$$(3.2.13) \quad \neg \circ_1 \perp \leftrightarrow \perp \circ_2 \neg,$$

$$(3.2.17) \quad \vdash \circ_2 \vdash \leftrightarrow \vdash \circ_1 \vdash \leftrightarrow \vdash \circ_1 \neg \leftrightarrow \vdash \circ_1 \perp.$$

$$(3.2.14) \quad \perp \circ_1 \neg \leftrightarrow \perp \circ_2 \vdash,$$

Proposition 3.25. *The ns operads \mathbf{Tr} and \mathbf{Trias} are isomorphic and the map*

$$(3.2.18) \quad \phi : \mathbf{Trias} \rightarrow \mathbf{Tr}$$

satisfying $\phi(\neg) = 10$, $\phi(\vdash) = 01$, and $\phi(\perp) = 11$ is an isomorphism.

Proof. By replacing each generator g of \mathbf{Trias} by $\phi(g)$, the generators 01 , 10 , and 11 of \mathbf{Tr} satisfy the same relations as the generators \neg , \vdash , and \perp of \mathbf{Trias} . Moreover, by Proposition 3.24, the Hilbert series of \mathbf{Tr} is

$$(3.2.19) \quad F(t) = \sum_{n \geq 1} (2^n - 1)t^n,$$

which also is the Hilbert series of \mathbf{Trias} . Then, there is no nontrivial relation of degree greater than two involving generators of \mathbf{Tr} . \square

Proposition 3.25 also shows that \mathbf{Tr} is a realization of the triassociative operad.

CONCLUDING REMARKS

We have presented the functorial construction \mathbf{T} producing an operad given a monoid. As we have seen, this construction is very rich from a combinatorial point of view since most of the obtained operads coming from usual monoids involve a wide range of combinatorial objects. There are various way to continue this work. Let us address here the main directions.

In the first place, it appears that we have somewhat neglected the fact that \mathbf{T} is a functor to operads and not only to ns ones. Indeed, except for the operads \mathbf{End} , \mathbf{PF} , \mathbf{PW} , and \mathbf{Per} , we only have regarded the obtained operads as ns ones. Computer experiments let us think that the dimensions of the operads \mathbf{PRT} , $\mathbf{FCat}^{(2)}$, \mathbf{Mozt} , \mathbf{DA} and \mathbf{SComp} seen as symmetric ones are, respectively, Sequences [A052882](#), [A050351](#), [A032181](#), [A101052](#), and [A001047](#) of [Slo]. Bijections between elements of these operads and combinatorial objects enumerated by these

sequences, together with presentations by generators and relations in this symmetric context, would be worthwhile.

Furthermore, we have considered T only in the category of sets, *i.e.*, it takes a monoid as input and constructs a set-operad as output. We can obviously extend the definition of T over the category of vector spaces. In that event, T would be a functor from the category of unital associative algebras to the category of operads in the category of vector spaces. It is thus natural to ask what operads T produces in this category.

Another line of research is the following. It is well-known that the Koszul dual (see [GK94] for Koszul duality of operads) of the operads *Dias* and *Trias* are respectively the dendriform *Dendr* [Lod01] and the tridendriform *TDendr* [LR04] operads. The tridendriform operad is a generalization of the dendriform operad and further generalizations were proposed, like the operads *Quad* [AL04] and *Ennea* [Ler04]. Since the operads *Di* and *Tr*, obtained from the T construction, are respectively isomorphic to the operads *Dias* and *Trias*, we can ask if there are generalizations of *Di* and *Tr* so that their Koszul duals provide generalizations of the operads *Dendr* and *TDendr*.

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